

```

EEEEEEEEEEEEEEEE  MMM      MMM  UUU      UUU  LLL      AAAAAAAAAA  TTTTTTTTTTTTTTTT
EEEEEEEEEEEEEEEE  MMM      MMM  UUU      UUU  LLL      AAAAAAAAAA  TTTTTTTTTTTTTTTT
EEEEEEEEEEEEEEEE  MMM      MMM  UUU      UUU  LLL      AAAAAAAAAA  TTTTTTTTTTTTTTTT
EEE               MMMMMM  MMMMMM  UUU      UUU  LLL      AAA          AAA  TTT
EEE               MMMMMM  MMMMMM  UUU      UUU  LLL      AAA          AAA  TTT
EEE               MMMMMM  MMMMMM  UUU      UUU  LLL      AAA          AAA  TTT
EEE               MMM      MMM      UUU      UUU  LLL      AAA          AAA  TTT
EEE               MMM      MMM      UUU      UUU  LLL      AAA          AAA  TTT
EEEEEEEEEEEEEEEE  MMM      MMM      UUU      UUU  LLL      AAA          AAA  TTT
EEEEEEEEEEEEEEEE  MMM      MMM      UUU      UUU  LLL      AAA          AAA  TTT
EEEEEEEEEEEEEEEE  MMM      MMM      UUU      UUU  LLL      AAA          AAA  TTT
EEE               MMM      MMM      UUU      UUU  LLL      AAA          AAA  TTT
EEE               MMM      MMM      UUU      UUU  LLL      AAA          AAA  TTT
EEE               MMM      MMM      UUU      UUU  LLL      AAA          AAA  TTT
EEE               MMM      MMM      UUU      UUU  LLL      AAA          AAA  TTT
EEE               MMM      MMM      UUU      UUU  LLL      AAA          AAA  TTT
EEEEEEEEEEEEEEEE  MMM      MMM      UUUUUUUUUUUUUUU  LLLLLLLLLLLLLLLLL  AAA          AAA  TTT
EEEEEEEEEEEEEEEE  MMM      MMM      UUUUUUUUUUUUUUU  LLLLLLLLLLLLLLLLL  AAA          AAA  TTT
EEEEEEEEEEEEEEEE  MMM      MMM      UUUUUUUUUUUUUUU  LLLLLLLLLLLLLLLLL  AAA          AAA  TTT

```

```

VV      VV      AAAAAA  XX      XX      AAAAAA  RRRRRRRR  IIIIII  TTTTTTTTTT  HH      HH
VV      VV      AAAAAA  XX      XX      AAAAAA  RRRRRRRR  IIIIII  TTTTTTTTTT  HH      HH
VV      VV      AA       AA      XX      XX      AA       AA      RR       RR  II       II  TT       TT  HH      HH
VV      VV      AA       AA      XX      XX      AA       AA      RR       RR  II       II  TT       TT  HH      HH
VV      VV      AA       AA      XX      XX      AA       AA      RR       RR  II       II  TT       TT  HH      HH
VV      VV      AA       AA      XX      XX      AA       AA      RRRRRRRR  II       II  TT       TT  HH      HH
VV      VV      AA       AA      XX      XX      AA       AA      RRRRRRRR  II       II  TT       TT  HHHHHHHHHH
VV      VV      AAAAAAAAAA  XX      XX      AAAAAAAAAA  RR       RR  II       II  TT       TT  HHHHHHHHHH
VV      VV      AAAAAAAAAA  XX      XX      AAAAAAAAAA  RR       RR  II       II  TT       TT  HH      HH
  VV      VV      AA       AA      XX      XX      AA       AA      RR       RR  II       II  TT       TT  HH      HH
  VV      VV      AA       AA      XX      XX      AA       AA      RR       RR  II       II  TT       TT  HH      HH
    VV      VV      AA       AA      XX      XX      AA       AA      RR       RR  II       II  TT       TT  HH      HH
      VV      VV      AA       AA      XX      XX      AA       AA      RR       RR  IIIIII  TT       TT  HH      HH
        VV      VV      AA       AA      XX      XX      AA       AA      RR       RR  IIIIII  TT       TT  HH      HH

```

```

....
....
....
....

```

```

LL      IIIIII  SSSSSSSS
LL      IIIIII  SSSSSSSS
LL      II      SS
LL      II      SS
LL      II      SS
LL      II      SS
LL      II      SSSSSS
LL      II      SSSSSS
LL      II      SS
LL      II      SS
LL      II      SS
LL      II      S
LLLLLLLLLL  IIIIII  SSSSSSSS
LLLLLLLLLL  IIIIII  SSSSSSSS

```

(2)	72	Declarations
(3)	133	VAX\$SUBP6 - Subtract Packed (6 Operand Format)
(4)	184	VAX\$ADDP6 - Add Packed (6 Operand Format)
(5)	239	VAX\$SUBP4 - Subtract Packed (4 Operand Format)
(6)	282	VAX\$ADDP4 - Add Packed (4 Operand Format)
(7)	336	ADDPx/SUBPx Common Initialization Code
(8)	457	ADD_PACKED - Add Two Packed Decimal Strings
(9)	615	ADD_PACKED_BYTE - Add Two Bytes Containing Decimal Digits
(10)	722	SUBTRACT_PACKED - Subtract Two Packed Decimal Strings
(11)	924	SUBTRACT_PACKED_BYTE - Subtract Two Bytes Containing Decimal Digits
(12)	1040	STORE_RESULT - Store Decimal String
(13)	1124.1	- CHECK_WRITE_ACCESS - Check Writability of Decimal String
(14)	1125	VAX\$MULP - Multiply Packed
(15)	1322	Common Exit Path for VAX\$MULP and VAX\$DIVP
(16)	1500	EXTEND_STRING_MULTIPLY - Multiply a String by a Number
(17)	1582	VAX\$DIVP - Divide Packed
(18)	1938	QUOTIENT_DIGIT - Get Next Digit in Quotient
(19)	2100	MULTIPLY_STRING - Multiply a String by a Number
(20)	2154	DECIMAL_OPERAND
(21)	2195	ARITH_ACCESS_VIO - Reflect an Access Violation
(22)	2245	Access Violation Handling for ADDPx and SUBPx
(23)	2319	Access Violation Handling for MULP and DIVP

:LJK0045  
-1

```

0000 1 .TITLE VAX$DECIMAL_ARITHMETIC - VAX-11 Packed Decimal Arithmetic Instructio
0000 2 .IDENT /V04-001/
0000 3
0000 4
0000 5 *****
0000 6 *
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0000 24 *
0000 25 *
0000 26 *****
0000 27
0000 28
0000 29 ++
0000 30 Facility:
0000 31
0000 32 VAX-11 Instruction Emulator
0000 33
0000 34 Abstract:
0000 35
0000 36 The routines in this module emulate the VAX-11 packed decimal
0000 37 instructions that perform arithmetic operations. These procedures can
0000 38 be a part of an emulator package or can be called directly after the
0000 39 input parameters have been loaded into the architectural registers.
0000 40
0000 41 The input parameters to these routines are the registers that
0000 42 contain the intermediate instruction state.
0000 43
0000 44 Environment:
0000 45
0000 46 These routines run at any access mode, at any IPL, and are AST
0000 47 reentrant.
0000 48
0000 49 Author:
0000 50
0000 51 Lawrence J. Kenah
0000 52
0000 53 Creation Date
0000 54
0000 55 19 October 1983
0000 56
0000 57 Modified by:

```

:LJK0045 0000 58 :  
:LJK0045 0000 .1 :  
:LJK0045 0000 .2 :  
:LJK0045 0000 .3 :  
:LJK0045 0000 .4 :  
:LJK0045 0000 .5 :  
0000 59 :  
0000 60 :  
0000 61 :  
0000 62 :  
0000 63 :  
0000 64 :  
0000 65 :  
0000 66 :  
0000 67 :  
0000 68 :  
0000 69 :  
0000 70 :--

V04-001 LJK0045 Lawrence J. Kenah 19-Sep-1984  
The result string in ADDP4 and SUBP4 must be probed for  
write access to insure that the instruction does not modify  
memory before trying to restart.

V01-003 LJK0037 Lawrence J. Kenah 17-Jul-1984  
Fix two minor bugs in exception handling code that caused  
MULP and DIVP tests to generate spurious access violations.

V01-002 LJK0024 Lawrence J. Kenah 21-Feb-1984  
Add code to handle access violations. Perform minor cleanup.  
Eliminate double use of R10 in MULP and DIVP.

V01-001 LJK0008 Lawrence J. Kenah 19-Oct-1983  
The emulation code for ADDP4, ADDP6, SUBP4, SUBP6, MULP and  
DIVP was moved into a separate module.

```

0000 72      .SUBTITLE      Declarations
0000 73
0000 74      ; Include files:
0000 75
0000 76      .NCCROSS        ; No cross reference for these
0000 77      .ENABLE        SUPPRESSION ; No symbol table entries either
0000 78
0000 79      ADDP4_DEF       ; Bit fields in ADDP4 registers
0000 80      ADDP6_DEF       ; Bit fields in ADDP6 registers
0000 81      DIVP_DEF        ; Bit fields in DIVP registers
0000 82      MULP_DEF        ; Bit fields in MULP registers
0000 83      SUBP4_DEF       ; Bit fields in SUBP4 registers
0000 84      SUBP6_DEF       ; Bit fields in SUBP6 registers
0000 85
0000 86      $PSLDEF        ; Define bit fields in PSL
0000 87      $SRMDEF        ; Define arithmetic trap codes
0000 88
0000 89      .DISABLE       SUPPRESSION ; Turn on symbol table again
0000 90      .CROSS          ; Cross reference is OK now
0000 91
0000 92      ; Symbol definitions
0000 93
0000 94      ; The architecture requires that R4 be zero on completion of an ADDP6 or
0000 95      ; SUBP6 instruction. If we did not have to worry about restarting
0000 96      ; instructions after an access violation, we could simply zero the saved
0000 97      ; R4 value on the code path that these two instructions have in common
0000 98      ; before they merge with the ADDP4 and SUBP4 routines. The ability to
0000 99      ; restart requires that we keep the original R4 around at least until no
0000 100     ; more access violations are possible. To accomplish this, we store the
0000 101     ; fact that R4 must be cleared on exit in R11, which also contains the
0000 102     ; evolving condition codes. We use bit 31, the compatibility mode bit
0000 103     ; because it is nearly impossible to enter the emulator with CM set.
0000 104
0000001F 0000 105     ADD_SUB_V_ZERO_R4 = PSLSV_CM
0000 106
0000 107     ; External declarations
0000 108
0000 109     .DISABLE            GLOBAL
0000 110
0000 111     .EXTERNAL -
0000 112     DECIMAL$BOUNDS_CHECK,-
0000 113     DECIMAL$BINARY_TO_PACKED_TABLE,-
0000 114     DECIMAL$PACKED_TO_BINARY_TABLE,-
0000 115     DECIMAL$STRIP_ZEROS_R0_RT,-
0000 116     DECIMAL$STRIP_ZEROS_R2_R3
0000 117
0000 118     .EXTERNAL -
0000 119     VAX$DECIMAL_EXIT,-
0000 120     VAX$DECIMAL_ACCVIO,-
0000 121     VAX$REFLECT_TRAP,-
0000 122     VAX$ROPRAND
0000 123
0000 124     ; PSECT Declarations:
0000 125
0000 126     .DEFAULT            DISPLACEMENT , WORD
0000 127
00000000 0000 128     .PSECT _VAX$CODE PIC,USR,CON,REL,LCL,SHR,EXE,RD,NOWRT,LONG

```

0000 129  
0000 130

BEGIN\_MARK\_POINT

-1

```

0000 133      .SUBTITLE      VAX$$SUBP6 - Subtract Packed (6 Operand Format)
0000 134      :
0000 135      : Functional Description:
0000 136      :
0000 137      :     In 6 operand format, the subtrahend string specified by the subtrahend
0000 138      :     length and subtrahend address operands is subtracted from the minuend
0000 139      :     string specified by the minuend length and minuend address operands.
0000 140      :     The difference string specified by the difference length and difference
0000 141      :     address operands is replaced by the result.
0000 142      :
0000 143      : Input Parameters:
0000 144      :
0000 145      :     R0 - sublen.rw      Number of digits in subtrahend string
0000 146      :     R1 - subaddr.ab     Address of subtrahend string
0000 147      :     R2 - minlen.rw     Number of digits in minuend string
0000 148      :     R3 - minaddr.ab    Address of minuend string
0000 149      :     R4 - diflen.rw   Number of digits in difference string
0000 150      :     R5 - difaddr.ab   Address of difference string
0000 151      :
0000 152      : Output Parameters:
0000 153      :
0000 154      :     R0 = 0
0000 155      :     R1 = Address of the byte containing the most significant digit of
0000 156      :           the subtrahend string
0000 157      :     R2 = 0
0000 158      :     R3 = Address of the byte containing the most significant digit of
0000 159      :           the minuend string
0000 160      :     R4 = 0
0000 161      :     R5 = Address of the byte containing the most significant digit of
0000 162      :           the string containing the difference
0000 163      :
0000 164      : Condition Codes:
0000 165      :
0000 166      :     N <- difference string LSS 0
0000 167      :     Z <- difference string EQL 0
0000 168      :     V <- decimal overflow
0000 169      :     C <- 0
0000 170      :
0000 171      : Register Usage:
0000 172      :
0000 173      :     This routine uses all of the general registers. The condition codes
0000 174      :     are recorded in R11 as the routine executes.
0000 175      :
0000 176      :
0000 177      : .ENABLE      LOCAL_BLOCK
0000 178      :
0000 179      VAX$$SUBP6::
OFFF 8F  BB 0000 180      PUSHR  #^M<R0,R1,R2,R3,R4,R5,R6,R7,R8,R9,R10,R11> ; Save the lot
59  01  9A 0004 181      MOVZBL #1,R9 ; Indicate that this is subtraction
06  11  0007 182      BRB 10$ ; Merge with ADDP6 code
  
```



```

0009 184 .SUBTITLE VAX$ADDP6 - Add Packed (6 Operand Format)
0009 185 :+
0009 186 : Functional Description:
0009 187 :
0009 188 : In 6 operand format, the addend 1 string specified by the addend 1
0009 189 : length and addend 1 address operands is added to the addend 2 string
0009 190 : specified by the addend 2 length and addend 2 address operands. The sum
0009 191 : string specified by the sum length and sum address operands is replaced
0009 192 : by the result.
0009 193 :
0009 194 : Input Parameters:
0009 195 :
0009 196 : R0 - add1len.rw Number of digits in first addend string
0009 197 : R1 - add1addr.ab Address of first addend string
0009 198 : R2 - add2len.rw Number of digits in second addend string
0009 199 : R3 - add2addr.ab Address of second addend string
0009 200 : R4 - sumlen.rw Number of digits in sum string
0009 201 : R5 - sumaddr.ab Address of sum string
0009 202 :
0009 203 : Output Parameters:
0009 204 :
0009 205 : R0 = 0
0009 206 : R1 = Address of the byte containing the most significant digit of
0009 207 : the first addend string
0009 208 : R2 = 0
0009 209 : R3 = Address of the byte containing the most significant digit of
0009 210 : the second addend string
0009 211 : R4 = 0
0009 212 : R5 = Address of the byte containing the most significant digit of
0009 213 : the string containing the sum
0009 214 :
0009 215 : Condition Codes:
0009 216 :
0009 217 : N <- sum string LSS 0
0009 218 : Z <- sum string EQL 0
0009 219 : V <- decimal overflow
0009 220 : C <- 0
0009 221 :
0009 222 : Register Usage:
0009 223 :
0009 224 : This routine uses all of the general registers. The condition codes
0009 225 : are recorded in R11 as the routine executes.
0009 226 :-
0009 227 :

```

```

0009 228 VAX$ADDP6::
OFFF 8F BB 0009 229 PUSHR #*M<R0,R1,R2,R3,R4,R5,R6,R7,R8,R9,R10,R11> ; Save the lot
59 D4 000D 230 CLR R9 ; This is addition
000F 231 10$: ROPRAND_CHECK R4 ; Insure that R4 is IEGU 31
5B DC 001A 232 MOVPSL R11 ; Get initial PSL
001C 233
001C 234 ; Indicate that the saved R4 must be cleared on the exit path
001C 235
:LJK0045 22 5B 1F E3 001C .1 BBCS #ADD_SUB_V_ZERO_R4,R11,25$ ; Set bit and join common code
:LJK0045 20 11 0020 .2 BRB 25$ ; In case we drop through BBCS

```

```

0022 239      .SUBTITLE      VAX$SUBP4 - Subtract Packed (4 Operand Format)
0022 240      :
0022 241      : Functional Description:
0022 242      :
0022 243      :     In 4 operand format, the subtrahend string specified by subtrahend
0022 244      :     length and subtrahend address operands is subtracted from the difference
0022 245      :     string specified by the difference length and difference address
0022 246      :     operands and the difference string is replaced by the result.
0022 247      :
0022 248      : Input Parameters:
0022 249      :
0022 250      :     R0 - sublen.rw      Number of digits in subtrahend string
0022 251      :     R1 - subaddr.ab   Address of subtrahend decimal string
0022 252      :     R2 - diflen.rw    Number of digits in difference string
0022 253      :     R3 - difaddr.ab   Address of difference decimal string
0022 254      :
0022 255      : Output Parameters:
0022 256      :
0022 257      :     R0 = 0
0022 258      :     R1 = Address of the byte containing the most significant digit of
0022 259      :           the subtrahend string
0022 260      :     R2 = 0
0022 261      :     R3 = Address of the byte containing the most significant digit of
0022 262      :           the string containing the difference
0022 263      :
0022 264      : Condition Codes:
0022 265      :
0022 266      :     N <- difference string LSS 0
0022 267      :     Z <- difference string EQL 0
0022 268      :     V <- decimal overflow
0022 269      :     C <- 0
0022 270      :
0022 271      : Register Usage:
0022 272      :
0022 273      :     This routine uses all of the general registers. The condition codes
0022 274      :     are recorded in R11 as the routine executes.
0022 275      : -
0022 276      :
0022 277      VAX$SUBP4::
OFFF 8F  BB 0022 278      PUSHR  #^M<R0,R1,R2,R3,R4,R5,R6,R7,R8,R9,R10,R11> ; Save the lot
59  01  9A 0026 279      MOVZBL #1,R9 ; Indicate that this is subtraction
0022 280      BRB 20$ ; Merge with ADDP4 code
  
```

```

002B 282 .SUBTITLE VAX$ADDP4 - Add Packed (4 Operand Format)
002B 283
002B 284 : Functional Description:
002B 285
002B 286 In 4 operand format, the addend string specified by the addend length
002B 287 and addend address operands is added to the sum string specified by the
002B 288 sum length and sum address operands and the sum string is replaced by
002B 289 the result.
002B 290
002B 291 : Input Parameters:
002B 292
002B 293 R0 - addlen.rw Number of digits in addend string
002B 294 R1 - addaddr.ab Address of addend decimal string
002B 295 R2 - sumlen.rw Number of digits in sum string
002B 296 R3 - sumaddr.ab Address of sum decimal string
002B 297
002B 298 : Output Parameters:
002B 299
002B 300 R0 = 0
002B 301 R1 = Address of the byte containing the most significant digit of
002B 302 the addend string
002B 303 R2 = 0
002B 304 R3 = Address of the byte containing the most significant digit of
002B 305 the string containing the sum
002B 306
002B 307 : Condition Codes:
002B 308
002B 309 N <- sum string LSS 0
002B 310 Z <- sum string EQL 0
002B 311 V <- decimal overflow
002B 312 C <- 0
002B 313
002B 314 : Register Usage:
002B 315
002B 316 This routine uses all of the general registers. The condition codes
002B 317 are recorded in R11 as the routine executes.
002B 318 :-
002B 319
:LJK0045 0264 30 002B .1 15$: BSBW CHECK_WRITE_ACCESS ; Perform rigorous access check
:LJK0045 38 11 002E .2 BRB 30$ ; String can be written after all
:LJK0045
0030 320 VAX$ADDP4::
0030 321 PUSHR #*M<R0,R1,R2,R3,R4,R5,R6,R7,R8,R9,R10,R11> ; Save the lot
0034 322 CLRL R9 ; This is addition
0036 323
0036 324 ; The output string, described by R4 and R5, will be the same as the input
0036 325 ; string for ADDP4 and SUBP4. It is necessary to explicitly clear R4<31:16>
0036 326 ; along this code path so MOVQ R2,R4 will not always work.
0036 327
0036 328 20$: MOVZWL R2,R4 ; Set output size equal to input size
0039 329 MOVL R3,R5 ; ... and ditto for string addresses
003C 330 MOVPSL R11 ; Get initial PSL
003E 331
003E 332 ; Indicate that the saved R4 will be restored on the common exit path
003E 333
:LJK0045 00 5B 1F E5 003E .1 BBCC #ADD_SUB_V_ZERO_R4,R11,25$ ; Clear bit and join common code

```

```

0042 336 .SUBTITLE ADDPx/SUBPx Common Initialization Code
0043 337
0044 338 : All four routines converge at this point and execute common initialization
0045 339 : code until a later decision is made to do addition or subtraction.
0046 340
0047 341 R4 - Number of digits in destination string
0048 342 R5 - Address of destination string
0049 343
0050 344 R9 - Indicates whether operation is addition or subtraction
0051 345 0 => addition
0052 346 1 => subtraction
0053 347
0054 348 R11<31> - Indicates whether this is a 4-operand or 6-operand instruction
0055 349 0 => 4-operand (restore saved R4 on exit)
0056 350 1 => 6-operand (set R4 to zero on exit)
0057 351 :-
0058 352
:LJK0045 04 00 04 F0 0042 .1 258: INSV #PSLSM_2,#0,#4,R11 ; Set Z-bit, clear the rest in saved PSW
-1 0046
0047 354 ESTABLISH HANDLER - ; Store address of access
0048 355 ARITH_ACCVIO ; violation handler
0049 356
0050 357 ROPRAND CHECK R2 ; Insure that R2 is LEQU 31
0051 358 MARK_POINT ADD SUB BSBW 0
FFA9' 30 0054 359 BSBW- DECIMAL$STRTP_ZEROS_R2_R3 ; Strip high order zeros from R2/R3
0055 360
0056 361 ROPRAND CHECK R0 ; Insure that R0 is LEQU 31
0057 362 MARK_POINT ADD SUB BSBW 0
FF9E' 30 005F 363 BSBW- DECIMAL$STRTP_ZEROS_R0_R1 ; Strip high order zeros from R0/R1
0058 364
:LJK0045 0062 364 : Perform the access check on the output string for the worst case, a string
:LJK0045 0062 .1 : large enough to accommodate 31 decimal digits. A detailed check using the
:LJK0045 0062 .2 : correct byte length of the output string is necessary only if this initial
:LJK0045 0062 .3 : probe fails. The more detailed check can be handled out of line.
:LJK0045 0062 .4
:LJK0045 0062 .5
:LJK0045 65 10 00 0D 0062 .6 PROBEW #0,#16,(R5) ; Can result string be written?
:LJK0045 0066 .7 BEQL 15$ ; Branch if no write access allowed
:LJK0045 0068 .8
0068 365 : Rather than totally confuse the already complicated logic dealing with
0069 366 : different length strings in the add or subtract loop, we will put the
0070 367 : result into an intermediate buffer on the stack. This buffer will be long
0071 368 : enough to handle the worst case so that the addition loop need only concern
0072 369 : itself with the lengths of the two input loops. The required length is 17
0073 370 : bytes to handle an addition with a carry out of the most significant byte.
0074 371 : We will allocate 20 bytes to maintain whatever alignment the stack has.
0075 372
:LJK0045 0068 .1 308: CLRQ -(SP) ; Set aside space for output string
-1 006A 374 CLRQ -(SP) ; Worst case string needs 16 bytes
006C 375 CLRL -(SP) ; Add slack for a CARRY
54 04 01 EF 006E 376 EXTZV #1,#4,R4,R8 ; Get byte count of destination string
0072
7E 55 58 C1 0073 377 ADDL3 R8,R5,-(SP) ; Save high address end of destination
55 18 AE 9E 0077 378 MOVAB 24(SP),R5 ; Point R5 one byte beyond buffer
007B 379
007B 380 : The number of minus signs will determine whether the real operation that we
007B 381 : perform is addition or subtraction. That is, two plus signs or two minus
007B 382 : signs will both result in addition, while a plus sign and a minus sign will

```

```

007B 383 ; result in subtraction. The addition and subtraction routines have their own
007B 384 ; methods for determining the correct sign of the result.
007B 385 ;
007B 386 ; For the purpose of counting minus signs, we treat subtraction as the
007B 387 ; addition of the negative of the input operand. That is, subtraction of a
007B 388 ; positive quantity causes the sign to be remembered as minus and counted as
007B 389 ; a minus sign while subtraction of a minus quantity stores a plus sign and
007B 390 ; counts nothing.
007B 391 ;
007B 392 ; On input to this code sequence, R9 distinguished addition from subtraction.
007B 393 ; On output, it contains either 0, 1, or 2, indicating the total number of
007B 394 ; minus signs, real or implied, that we counted.
007B 395 ;
50 04 01 EF 007B 396      EXTZV  #1,#4,R0,R6      ; Get byte count for first input string
                                007F
                                51 56 CO 0080 397      ADDL   R6,R1      ; Point R1 to byte containing sign
                                0083 398      MARK POINT  ADD SUB 24
61  F0 8F 8B 0083 399      BICB3  #*B11110000,(R1),R6  ; R6 contains the sign 'digit'
                                0087
                                10 59 E8 0088 400      BLBS   R9,35$     ; Use second CASE if subtraction
008B 401
008B 402 ; This case statement is used for addition
008B 403
008B 404      CASE   R6,TYPE=B,LIMIT=#10,<-  ; Dispatch on sign digit
008B 405      50$,-  ; 10 => sign is '+'
008B 406      40$,-  ; 11 => sign is '-'
008B 407      50$,-  ; 12 => sign is '+'
008B 408      40$,-  ; 13 => sign is '-'
008B 409      50$,-  ; 14 => sign is '+'
008B 410      50$,-  ; 15 => sign is '+'
008B 411      >
009B 412
009B 413 ; This case statement is used for subtraction
009B 414
009B 415 35$:  CASE   R6,TYPE=B,LIMIT=#10,<-  ; Dispatch on sign digit
009B 416      40$,-  ; 10 => treat sign as '-'
009B 417      50$,-  ; 11 => treat sign as '+'
009B 418      40$,-  ; 12 => treat sign as '-'
009B 419      50$,-  ; 13 => treat sign as '+'
009B 420      40$,-  ; 14 => treat sign as '-'
009B 421      40$,-  ; 15 => treat sign as '-'
009B 422      >
00AB 423
59 01 D0 00AB 424 40$:  MOVL   #1,R9      ; Count a minus sign
56 0D 9A 00AE 425      MOV7BL #13,R6     ; The preferred minus sign is 13
                                05 11 00B1 426      BRB    60$      ; Now check second input sign
                                00B3 427
                                56 59 D4 00B3 428 50$:  CLRL   R9      ; No real minus signs so far
                                0C 9A 00B5 429      MOVZBL #12,R6     ; The preferred minus sign is 12
                                00B8 430
52 04 01 EF 00B8 431 60$:  EXTZV  #1,#4,R2,R7  ; Get byte count for second input string
                                00BC
                                53 57 CO 00BD 432      ADDL   R7,R3      ; Point R3 to byte containing sign
                                00C0 433      MARK POINT  ADD SUB 24
63  F0 8F 8B 00C0 434      BICB3  #*B11110000,(R3),R7  ; R7 contains the sign 'digit'
                                00C4
                                00C5 435

```

```
00C5 436      CASE R7,TYPE=B,LIMIT=#10,<- ; Dispatch on sign digit
00C5 437      80$,- ; 10 => sign is '+'
00C5 438      70$,- ; 11 => sign is '-'
00C5 439      80$,- ; 12 => sign is '+'
00C5 440      70$,- ; 13 => sign is '-'
00C5 441      80$,- ; 14 => sign is '+'
00C5 442      80$,- ; 15 => sign is '+'
00C5 443      >
00D5 444
57 59 D6 00D5 445 70$: INCL R9 ; Remember that sign was minus
03 0D 9A 00D7 446 MOVZBL #13,R7 ; The preferred minus sign is 13
03 03 11 00DA 447 BRB 90$ ; Now check second input sign
00DC 448
57 0C 9A 00DC 449 80$: MOVZBL #12,R7 ; The preferred minus sign is 12
00DF 450
03 59 E9 00DF 451 90$: BLBC R9,ADD_PACKED ; Even parity indicates addition
00E2 452
00B3 31 00E2 453 BRW SUBTRACT_PACKED ; Odd parity calls for subtraction
00E5 454
00E5 455 .DISABLE LOCAL_BLOCK
```

00E5 457  
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00E5 513

.SUBTITLE ADD\_PACKED - Add Two Packed Decimal Strings

Functional Description:

This routine adds two packed decimal strings whose descriptors are passed as input parameters and places their sum into another (perhaps identical) packed decimal string.

At the present time, the result is placed into a 16-byte storage area while the sum is being evaluated. This drastically reduces the number of different cases that must be dealt with as each pair of bytes in the two input strings is added.

The signs of the two input strings have already been dealt with so this routine performs addition in all cases, even if the original entry was at SUBP4 or SUBP6. The cases that arrive in this routine are as follows.

	R2/R3	R0/R1	result
R2/R3 + R0/R1	plus	plus	plus
R2/R3 + R0/R1	minus	minus	minus
R2/R3 - R0/R1	minus	plus	minus
R2/R3 - R0/R1	plus	minus	plus

Note that the correct choice of sign in all four cases is the sign of the second input string, the one described by R2 and R3.

Input Parameters:

- R0<4:0> - Number of digits in first input decimal string
- R1 - Address of least significant digit of first input decimal string (the byte containing the sign)
- R2<4:0> - Number of digits in second input decimal string
- R3 - Address of least significant digit of second input decimal string (the byte containing the sign)
- R4<4:0> - Number of digits in output decimal string
- R5 - Address of one byte beyond least significant digit of intermediate string stored on the stack
- R6<3:0> - Sign of first input string in preferred form
- R7<3:0> - Sign of second input string in preferred form







```

016A 615      .SUBTITLE      ADD_PACKED_BYTE - Add Two Bytes Containing Decimal Digits
016A 616      :
016A 617      : * Functional Description:
016A 618      :
016A 619      : This routine adds together two bytes containing decimal digits and
016A 620      : produces a byte containing the sum that is stored in the output
016A 621      : string. Each of the input bytes is converted to a binary number
016A 622      : (with a table-driven conversion), the two numbers are added, and
016A 623      : the sum is converted back to two decimal digits stored in a byte.
016A 624      :
016A 625      : This routine makes no provisions for bytes that contain illegal
016A 626      : decimal digits. We are using the UNPREDICTABLE statement in the
016A 627      : architectural description of the decimal instructions to its fullest.
016A 628      :
016A 629      : The bytes that contain a pair of packed decimal digits can either
016A 630      : exist in packed decimal strings located by R1 and R3 or they can
016A 631      : be stored directly in registers. In the former case, the digits must
016A 632      : be extracted from registers before they can be used in later operations
016A 633      : because the sum will be used as an index register.
016A 634      :
016A 635      : For entry at ADD_PACKED_BYTE_STRING:
016A 636      :
016A 637      : Input Parameters:
016A 638      :
016A 639      :     R1 - Address one byte beyond first byte that is to be added
016A 640      :     R3 - Address one byte beyond second byte that is to be added
016A 641      :     R5 - Address one byte beyond location to store sum
016A 642      :
016A 643      :     R8 - Carry from previous byte (R8 is either 0 or 1)
016A 644      :
016A 645      : Implicit Input:
016A 646      :
016A 647      :     R6 - Scratch
016A 648      :     R7 - Scratch
016A 649      :
016A 650      : Output Parameters:
016A 651      :
016A 652      :     R1 - Decreased by one to point to current byte in first input string
016A 653      :     R3 - Decreased by one to point to current byte in second input string
016A 654      :     R5 - Decreased by one to point to current byte in output string
016A 655      :
016A 656      :     R8 - Either 0 or 1, reflecting whether this most recent ADD resulted
016A 657      :           in a CARRY to the next byte.
016A 658      :
016A 659      : For entry at ADD_PACKED_BYTE_R6_R7:
016A 660      :
016A 661      : Input Parameters:
016A 662      :
016A 663      :     R6 - First byte containing decimal digit pair
016A 664      :     R7 - Second byte containing decimal digit pair
016A 665      :
016A 666      :     R5 - Address one byte beyond location to store sum
016A 667      :
016A 668      :     R8 - Carry from previous byte (R8 is either 0 or 1)
016A 669      :
016A 670      : Output Parameters:
016A 671      :
  
```

```

016A 672 : R5 - Decreased by one to point to current byte in output string
016A 673 :
016A 674 : R8 - Either 0 or 1, reflecting whether this most recent ADD resulted
016A 675 : in a CARRY to the next byte.
016A 676 :
016A 677 : Side Effects:
016A 678 : R6 and R7 are modified by this routine
016A 679 :
016A 680 : R0, R2, R4, and R9 (and, of course, R10 and R11) are preserved
016A 681 : by this routine
016A 682 :
016A 683 : Assumptions:
016A 684 : This routine makes two important assumptions.
016A 685 :
016A 686 : 1. If both of the input bytes contain only legal decimal digits, then
016A 687 : it is only necessary to subtract 100 at most once to put all
016A 688 : possible sums in the range 0..99. That is,
016A 689 :
016A 690 : 99 + 99 + 1 = 199 LSS 200
016A 691 :
016A 692 : 2. The result will be checked in some way to determine whether the
016A 693 : result is nonzero so that the Z-bit can have its correct setting.
016A 694 :
016A 695 :
016A 696 :-
016A 697 :-
016A 698 ADD_PACKED_BYTE_STRING:
016A 699
56 71 9A 016A 700 MARK POINT ADD_SUB_BSBW_24
016A 701 MOVZBL -(R1),R6 ; Get byte from first string
016D 702 MARK POINT ADD_SUB_BSBW_24
57 73 9A 016D 703 MOVZBL -(R3),R7 ; Get byte from second string
0170 704
0170 705 VAX$ADD_PACKED_BYTE_R6_R7:: ; ASHP also uses this routine
0170 706 ADD_PACKED_BYTE_R6_R7:-
56 0000'CF46 90 0170 707 MOVB DECIMAL$PACKED_TO_BINARY_TABLE[R6],-
0176 708 R6 ; Convert digits to binary
57 0000'CF47 90 0176 709 MOVB DECIMAL$PACKED_TO_BINARY_TABLE[R7],-
017C 710 R7 ; Convert digits to binary
57 56 80 017C 711 ADDB R6,R7 ; Form their sum
57 58 80 017F 712 ADDB R8,R7 ; Add CARRY from last step
63 8F 57 94 0182 713 CLRB R8 ; Assume no CARRY this time
58 07 18 0184 714 CMPB R7,#99 ; Check for CARRY
58 01 90 0188 715 BLEQU 10$ ; Branch if within bounds
75 57 64 8F 82 018D 716 MOVB #1,R8 ; Propagate CARRY to next step
0000'CF47 90 0191 717 SUBB #100,R7 ; Put R7 into interval 0..99
0197 718 10$: MOVB DECIMAL$BINARY_TO_PACKED_TABLE[R7],-
0197 719 -(R5) ; Store converted sum byte
05 0197 720 RSB

```

0198 722 .SUBTITLE SUBTRACT\_PACKED - Subtract Two Packed Decimal Strings  
0198 723

0198 724 : Functional Description:  
0198 725

0198 726 : This routine takes two packed decimal strings whose descriptors  
0198 727 : are passed as input parameters, subtracts one string from the  
0198 728 : other, and places their sum into another (perhaps identical)  
0198 729 : packed decimal string.  
0198 730

0198 731 : At the present time, the result is placed into a 16-byte storage  
0198 732 : area while the difference is being evaluated. This drastically reduces  
0198 733 : the number of different cases that must be dealt with as each  
0198 734 : pair of bytes in the two input strings is added.  
0198 735

0198 736 : The signs of the two input strings have already been dealt with so  
0198 737 : this routine performs subtraction in all cases, even if the original  
0198 738 : entry was at ADDP4 or ADDP6.  
0198 739

0198 740 : Input Parameters:  
0198 741

0198 742 : R0<4:0> - Number of digits in first input decimal string  
0198 743 : R1 - Address of least significant digit of first input  
0198 744 : decimal string (the byte containing the sign)  
0198 745

0198 746 : R2<4:0> - Number of digits in second input decimal string  
0198 747 : R3 - Address of least significant digit of second input  
0198 748 : decimal string (the byte containing the sign)  
0198 749

0198 750 : R4<4:0> - Number of digits in output decimal string  
0198 751 : R5 - Address of one byte beyond least significant digit of  
0198 752 : intermediate string stored on the stack  
0198 753

0198 754 : R6<3:0> - Sign of first input string in preferred form  
0198 755 : R7<3:0> - Sign of second input string in preferred form  
0198 756

0198 757 : R11 - Saved PSL (Z-bit is set, other condition codes are clear)  
0198 758 : (SP) - Saved R5, address of least significant digit of ultimate  
0198 759 : destination string.  
0198 760 : 4(SP) - Beginning of 20-byte buffer to hold intermediate result  
0198 761

0198 762 : Output Parameters:  
0198 763

0198 764 : The particular input operation (ADDPx or SUBPx) is completed in  
0198 765 : this routine. See the routine headers for the four routines that  
0198 766 : request addition or subtraction for a list of output parameters  
0198 767 : from this routine.  
0198 768  
0198 769

0198 770 : Algorithm for Choice of Sign:  
0198 771

0198 772 : The choice of sign for the output string is not nearly so  
0198 773 : straightforward as it is in the case of addition. One approach that is  
0198 774 : often taken is to make a reasonable guess at the sign of the result.  
0198 775 : If the final subtraction causes a BORROW, then the choice was incorrect.  
0198 776 : The sign must be changed and the result must be replaced by its tens  
0198 777 : complement.  
0198 778

0198 779 : This routine does not guess. Instead, it chooses the input string of  
0198 780 : the larger absolute magnitude as the minuend for this internal  
0198 781 : routine and chooses its sign as the sign of the result.  
0198 782 : This algorithm is actually more efficient than the reasonable  
0198 783 : guess method and is probably better than a guess method that is never  
0198 784 : wrong. All complete bytes that are processed in the sign evaluation  
0198 785 : preprocessing loop are eliminated from consideration in the  
0198 786 : subtraction loop, which has a higher cost per byte.

0198 787 :  
0198 788 : The actual algorithm is as follows. (Note that both input strings have  
0198 789 : already had leading zeros stripped so their lengths reflect  
0198 790 : significant digits.)

- 0198 791 : 1. If the two strings have unequal lengths, then choose the sign of  
0198 792 : the string that has the longer length.
- 0198 793 : 2. For strings of equal length, choose the sign of the string whose  
0198 794 : most significant byte is larger in magnitude.
- 0198 795 : 3. If the most significant bytes test equal, then decrease the  
0198 796 : lengths of each string by one byte, drop the previous most  
0198 797 : significant bytes, and go back to step 2.
- 0198 798 : 4. If the two strings test equal, it is not necessary to do any  
0198 799 : subtraction. The result is identically zero.

0198 800 : Note that the key to this routine's efficiency is that high order  
0198 801 : bytes that test equal in this loop are dropped from consideration in  
0198 802 : the more complicated subtraction loop.  
0198 803 :  
0198 804 :  
0198 805 :  
0198 806 :  
0198 807 :  
0198 808 :  
0198 809 :-

```

0198 810 SUBTRACT PACKED:
50 04 01 EF 0198 811 EXTZV #1,#4,R0,R0 ; Convert digit count to byte count
      50 019C
52 04 01 EF 019D 812 EXTZV #1,#4,R2,R2 ; Do it for both strings
      52 52 01A1
      50 D1 01A2 813 CMPL R0,R2 ; We want to compare the byte counts
      3C 1F 01A5 814 BLSSU 40$ ; R0/R1 represent the smaller string
      2A 1A 01A7 815 BGTRU 30$ ; R2/R3 represent the smaller string
      01A9 816
      01A9 817 The two input strings have an equal number of bytes. Compare magnitudes to
      01A9 818 : determine which string is really larger. If the two strings test equal, then
      01A9 819 : skip the entire subtraction loop.
      01A9 820
58 51 50 C3 01A9 821 SUBL3 R0,R1,R8 ; Point R8 to low address end of R0/R1
59 53 52 C3 01AD 822 SUBL3 R2,R3,R9 ; Point R9 to low address end of R2/R3
      50 D5 01B1 823 TSTL R0 ; See if both strings have zero bytes
      0C 13 01B3 824 BEQL 20$ ; Still need to check low order digit
      01B5 825
      01B5 826 MARK_POINT ADD_SUB_24
      89 88 91 01B5 827 10$: CMPB (R8)+,(R9)+ ; Compare most significant bytes
      29 1F 01B8 828 BLSSU 40$ ; R0/R1 represent the smaller string
      17 1A 01BA 829 BGTRU 30$ ; R2/R3 represent the smaller string
      52 D7 01BC 830 DECL R2 ; Keep R2 in step with R0
      F4 50 F5 01BE 831 SOBGTR R0,10$ ; ... which gets decremented here
      01C1 832
      01C1 833 ; At this point, we have reduced both input strings to single bytes that

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```

01C1 834 ; contain a sign 'digit' and may contain a digit in the high order nibble
01C1 835 ; if the original digit counts were nonzero.
01C1 836
01C1 837
58 68 OF 8B 01C1 838 20$: MARK POINT ADD SUB 24
01C1 839 BICB3 #^B00001111,(R0),R8 ; Look only at digit, ignoring sign
59 69 OF 8B 01C5 840 MARK POINT ADD SUB 24
59 59 58 91 01C9 841 BICB3 #^B00001111,(R9),R9 ; Get the digit from the other string
15 1F 01CC 842 CMPB R8,R9 ; Compare these digits
03 1A 01CE 843 BLSSU 40$ ; R0/hi represent the smaller string
01D0 844 BGTRU 30$ ; R2/R3 represent the smaller string
01D0 845 ; The two strings have identical magnitudes. Enter the end processing code
01D0 846 ; with the intermediate result unchanged (that is, zero).
01D0 847
FF6A 31 01D0 848 BRW ADD_SUBTRACT_EXIT ; Join the common completion code
01D3 849
01D3 850 ; The string described by R0 and R1 has the larger magnitude. Choose its sign.
01D3 851 ; Then swap the two string descriptors so that the main subtraction loops
01D3 852 ; always have R2 and R3 describing the larger string. Note that the use of
01D3 853 ; R6 and R7 as scratch leaves R7<31:8> in an UNPREDICTABLE state.
01D3 854
59 56 90 01D3 855 30$: MOVB R6,R9 ; Load preferred sign into R9
56 50 7D 01D6 856 MOVQ R0,R6 ; Save the longer
50 52 7D 01D9 857 MOVQ R2,R0 ; Store the shorter on R0 and R1
52 56 7D 01DC 858 MOVQ R6,R2 ; ... and store the longer in R2 and R3
57 57 D4 01DF 859 CLRL R7 ; Insure that R7<31:8> is zero
03 11 01E1 860 BRB 50$ ; Continue along common code path
01E3 861
01E3 862 ; The string described by P2 and R3 has the larger magnitude. Choose its sign.
01E3 863
59 57 90 01E3 864 40$: MOVB R7,R9 ; Load preferred sign into R9
01E6 865
52 50 C2 01E6 866 50$: SUBL R0,R2 ; Make R2 a difference (R2 GEQU 0)
03 59 E9 01E9 867 BLBC R9,60$ ; Check if sign is negative
58 08 88 01EC 868 BISB #PSLSM_N,R11 ; ... so the saved N-bit can be set
01EF 869
01EF 870
56 61 OF 8B 01EF 871 60$: MARK POINT ADD SUB 24
01F3 872 BICB3 #^B00001111,(R1),R6 ; Get least significant digit to R6
57 63 OF 8B 01F3 873 MARK POINT ADD SUB 24
58 D4 01F7 874 BICB3 #^B00001111,(R3),R7 ; Get least significant digit to R7
0032 30 01F9 875 CLRL R8 ; Start subtracting with BORROW off
01FC 876 BSBW SUB_PACKED_BYTE_R6_R7 ; Subtract the two low order digits
01FC 877 ; R0 contains the number of bytes remaining in the smaller string
01FC 878 ; R2 contains the difference in bytes between the two input strings
01FC 879
50 D5 01FC 880 TSTL R0 ; Does smaller string have any room?
06 13 01FE 881 BEQL 80$ ; Skip loop if no room at all
0200 882
0025 30 0200 883 70$: BSBW SUB_PACKED_BYTE_STRING ; Subtract the next two bytes
FA 50 F5 0203 884 SOBGTR R0,70$ ; Check for end of loop
0206 885
52 D5 0206 886 80$: TSTL R2 ; Does one of the strings have more?
16 13 0208 887 BEQL 110$ ; Skip next loops if all done
020A 888
OD 58 E9 020A 889 90$: BLBC R8,100$ ; Life is simple if BORROW clear
020D 890

```

```

56 D4 220D 891 CLR R6 ; Otherwise, BORROW must propagate
      020E 892 MARK POINT ADD_SUB_24
57 73 9A 020F 893 MOVZBL -(R3),R7 ; So subtract BORROW from single string
  0019 30 0212 894 BSBW SUB_PACKED_BITS_R6_R7 ; Use the special entry point
  F2 52 F5 0215 895 SOBGTR R2,90$ ; Check for this string exhausted
  06 11 0218 896 BRB 110$ ; Join common completion code
      021A 897
      021A 898 MARK POINT ADD_SUB_24
75 73 90 021A 899 100$: MOVB -(R3),-(R5) ; Simply move src to dst if no BORROW
  FA 52 F5 021D 900 SOBGTR R2,100$ ; ... until we're all done
      0220 901
      0220 902 110$:
      0220 903
      0220 904 ::: ***** BEGIN TEMP *****
      0220 905 :::
      0220 906 ::: THE FOLLOWING HALT INSTRUCTION SHOULD BE REPLACED WITH THE CORRECT
      0220 907 ::: ABORT CODE.
      0220 908 :::
      0220 909 ::: THE HALT IS SIMILAR TO THE
      0220 910 :::
      0220 911 ::: MICROCODE CANNOT GET HERE
      0220 912 :::
      0220 913 ::: ERRORS THAT OTHER IMPLEMENTATIONS USE.
      0220 914 :::
58 D5 0220 915 tstl r8 ; If BORROW is set here, we blew it
  01 13 0222 916 beql 120$ ; Branch out if OK
      00 0224 917 halt ; This will cause an OPCDEC exception
      0225 918 120$:
      0225 919 :::
      0225 920 ::: ***** END TEMP *****
      0225 921
FF15 31 0225 922 BRW ADD_SUBTRACT_EXIT ; Join common completion code

```

```

0228 924      .SUBTITLE      SUB_PACKED_BYTE - Subtract Two Bytes Containing Decimal Digi
0228 925      :
0228 926      : Functional Description:
0228 927      :
0228 928      : This routine takes as input two bytes containing decimal digits and
0228 929      : produces a byte containing their difference. This result is stored in
0228 930      : the output string. Each of the input bytes is converted to a binary
0228 931      : number (with a table-driven conversion), the first number is
0228 932      : subtracted from the second, and the difference is converted back to
0228 933      : two decimal digits stored in a byte.
0228 934      :
0228 935      : This routine makes no provisions for bytes that contain illegal
0228 936      : decimal digits. We are using the UNPREDICTABLE statement in the
0228 937      : architectural description of the decimal instructions to its fullest.
0228 938      :
0228 939      : The bytes that contain a pair of packed decimal digits can either
0228 940      : exist in packed decimal strings located by R1 and R3 or they can
0228 941      : be stored directly in registers. In the former case, the digits must
0228 942      : be extracted from registers before they can be used in later operations
0228 943      : because the difference will be used as an index register.
0228 944      :
0228 945      : For entry at SUB_PACKED_BYTE_STRING:
0228 946      :
0228 947      : Input Parameters:
0228 948      :
0228 949      :     R1 - Address one byte beyond byte containing subtrahend
0228 950      :     R3 - Address one byte beyond byte containing minuend
0228 951      :     R5 - Address one byte beyond location to store difference
0228 952      :
0228 953      :     R8 - BORROW from previous byte (R8 is either 0 or 1)
0228 954      :
0228 955      : Implicit Input:
0228 956      :
0228 957      :     R6 - Scratch
0228 958      :     R7 - Scratch
0228 959      :
0228 960      : Output Parameters:
0228 961      :
0228 962      :     R1 - Decreased by one to point to current byte
0228 963      :         in subtrahend string
0228 964      :     R3 - Decreased by one to point to current byte
0228 965      :         in minuend string
0228 966      :     R5 - Decreased by one to point to current byte
0228 967      :         in difference string
0228 968      :
0228 969      :     R8 - Either 0 or 1, reflecting whether this most recent
0228 970      :         subtraction resulted in a BORROW from the next byte.
0228 971      :
0228 972      : For entry at SUB_PACKED_BYTE_R6_R7:
0228 973      :
0228 974      : Input Parameters:
0228 975      :
0228 976      :     R6<7:0> - Byte containing decimal digit pair for subtrahend
0228 977      :     R6<31:8> - MBZ
0228 978      :     R7<7:0> - Byte containing decimal digit pair for minuend
0228 979      :     R7<31:8> - MBZ
0228 980      :

```



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0228 981 : R5 - Address one byte beyond location to store difference
0228 982 :
0228 983 : R8 - BORROW from subtraction of previous byte
0228 984 : (R8 is either 0 or 1)
0228 985 :
0228 986 : Output Parameters:
0228 987 :
0228 988 : R5 - Decreased by one to point to current byte
0228 989 : in difference string
0228 990 :
0228 991 : R8 - Either 0 or 1, reflecting whether this most recent
0228 992 : subtraction resulted in a BORROW from the next byte.
0228 993 :
0228 994 : Side Effects:
0228 995 :
0228 996 : R6 and R7 are modified by this routine
0228 997 :
0228 998 : R0, R2, R4, and R9 (and, of course, R10 and R11) are preserved
0228 999 : by this routine
0228 1000 :
0228 1001 : Assumptions:
0228 1002 :
0228 1003 : This routine makes two important assumptions.
0228 1004 :
0228 1005 : 1. If both of the input bytes contain only legal decimal digits, then
0228 1006 : it is only necessary to add 100 at most once to put all
0228 1007 : possible differences in the range 0..99. That is,
0228 1008 :
0228 1009 : 0 - 99 - 1 = -100
0228 1010 :
0228 1011 : 2. The result will be checked in some way to determine whether the
0228 1012 : result is nonzero so that the Z-bit can have its correct setting.
0228 1013 :
0228 1014 :
0228 1015 : SUB_PACKED_BYTE_STRING:
0228 1016 :
0228 1017 : MARK POINT ADD_SUB_BSBW_24
56 71 9A 0228 1018 : MOVZBL -(R1),R6 : Get byte from first string
0228 1019 : MARK POINT ADD_SUB_BSBW_24
57 73 9A 0228 1020 : MOVZBL -(R3),R7 : Get byte from second string
022E 1021 :
022E 1022 : SUB_PACKED_BYTE_R6_R7:
56 0000'CF46 90 022E 1023 : MOVB DECIMALSPACKED_TO_BINARY_TABLE[R6],-
0234 1024 : R6 : Convert digits to binary
57 0000'CF47 90 0234 1025 : MOVB DECIMALSPACKED_TO_BINARY_TABLE[R7],-
023A 1026 : R7 : Convert digits to binary
57 56 82 023A 1027 : SUBB R6,R7 : Form their difference
57 58 82 023D 1028 : SUBB R8,R7 : Include BORROW from last step
04 19 0240 1029 : BLSS 10$ : Branch if need to BORROW
58 94 0242 1030 : CLRB R8 : No BORROW next time
07 11 0244 1031 : BRB 20$ : Join common exit code
0246 1032 :
57 64 8F 80 0246 1033 10$: : ADDB #100,R7 : Put R7 into interval 0..99
58 01 90 024A 1034 : MOVB #1,R8 : Propagate BORROW to next step
024D 1035 :
75 0000'CF47 90 024D 1036 20$: : MOVB DECIMALSBINARY_TO_PACKED_TABLE[R7],-
0253 1037 : -(R5) : Store converted sum byte

```

VAX\$DECIMAL\_ARITHMETIC  
V04-001

N 8  
- VAX-11 Packed Decimal Arithmetic Instr 8-JAN-1985 17:27:01 VAX/VMS Macro V04-00 Page 23  
SUB\_PACKED\_BYTE - Subtract Two Bytes Con 5-SEP-1994 00:44:34 [EMULAT.BUGSRC]VAXARITH.MAR;1 (11)

05 0253 1038 RSB

```

0254 1040 .SUBTITLE STORE_RESULT - Store Decimal String
0254 1041 :
0254 1042 : Functional Description:
0254 1043 :
0254 1044 : This routine takes a packed decimal string that typically contains
0254 1045 : the result of an arithmetic operation and stores it in another
0254 1046 : decimal string whose descriptor is specified as an input parameter
0254 1047 : to the original arithmetic operation.
0254 1048 :
0254 1049 : The string is stored from the high address end (least significant
0254 1050 : digits) to the low address end (most significant digits). This order
0254 1051 : allows all of the special cases to be handled in the simplest fashion.
0254 1052 :
0254 1053 : Input Parameters:
0254 1054 :
0254 1055 : R1 - Address one byte beyond high address end of input string
0254 1056 : (Note that this string must be at least 17 bytes long.)
0254 1057 :
0254 1058 : R4<4:0> - Number of digits in ultimate destination
0254 1059 : R5 - Address one byte beyond destination string
0254 1060 :
0254 1061 : R11 - Contains saved condition codes
0254 1062 :
0254 1063 : Implicit Input:
0254 1064 :
0254 1065 : The input string must be at least 17 bytes long to contain a potential
0254 1066 : carry out of the highest digit when doing an add of two large numbers.
0254 1067 : This carry out of the last byte will be detected and reported as a
0254 1068 : decimal overflow, either as an exception or simply by setting the V-bit.
0254 1069 :
0254 1070 : The least significant digit (highest addressed byte) cannot contain a
0254 1071 : sign digit because that would cause the Z-bit to be incorrectly cleared.
0254 1072 :
0254 1073 : Output Parameters:
0254 1074 :
0254 1075 : R11<PSLSV_Z> - Cleared if a nonzero digit is stored in output string
0254 1076 : R11<PSLSV_V> - Set if a nonzero digit is detected after the output
0254 1077 : string is exhausted
0254 1078 :
0254 1079 : A portion of the result (dictated by the size of R4 on input) is
0254 1080 : moved to the destination string.
0254 1081 :-
0254 1082 :
0254 1083 STORE_RESULT:
0254 1084 INCL R4 ; Want number of "complete" bytes in
54 FF 8F 78 0256 1085 ASHL #-1,R4,R0 ; output string
025A 0B 13 025B 1086 BEQL 30$ ; Skip first loop if none
025D 1087
025D 1088 MARK_POINT ADD_SUB_BSBW_24
75 71 90 025D 1089 10$: MOVB -(R1),-(R5) ; Move the next complete byte
0260 1090 BEQL 20$ ; Check whether to clear Z-bit
5B 04 8A 0262 1091 BICB #PSLSM_Z,R11 ; Clear Z-bit if nonzero
F5 50 F5 0265 1092 20$: SOBGTR R0,10$ ; Keep going?
0268 1093
10 54 E9 0268 1094 30$: BLBC R4,50$ ; Was original R4 odd? Branch if yes
026B 1095 MARK_POINT ADD_SUB_BSBW_24

```

```

71  F0 8F 8B 026B 1096      BICB3  #^B11110000,-(R1),-(R5) ; If R4 was even, store half a byte
      75 026F
      03 13 0270 1097      BEQL   40$ ; Need to check for zero here, too
5B  04 8A 0272 1098      BICB  #PSLSM_Z,R11 ; Clear Z-bit if nonzero
      0275 1099      MARK_POINT ADD_SUB_BSBW_24
61  F0 8F 93 0275 1100 40$: BITB  #^B11110000,(R1) ; If high order nibble is nonzero,
      13 12 0279 1101      BNEQ   70$ ; ... then overflow has occurred
      027B 1102
      027B 1103 ; The entire destination has been stored. We must now check whether any of
      027B 1104 ; the remaining input string is nonzero and set the V-bit if nonzero is
      027B 1105 ; detected. Note that at least one byte of the output string has been examined
      027B 1106 ; in all cases already. This makes the next byte count calculation correct.
      027B 1107
54  04 54 D7 027B 1108 50$: DECL   R4 ; Restore R4 to its original self
      01 EF 027D 1109      EXTZV  #1,#4,R4,R0 ; Extract a byte count
      50 10 50 83 0282 1110      SUBB3  R0,#16,R0 ; Loop count is 16 minus byte count
      0286 1111
      0286 1112 ; Note that the loop count can never be zero because we are testing a 17-byte
      0286 1113 ; string and the largest output string can be 16 bytes long.
      0286 1114
      0286 1115
      71 95 0286 1116 60$: MARK_POINT ADD_SUB_BSBW_24
      04 12 0288 1117      TSTB  -(R1) ; Check next byte for nonzero
      F9 50 F5 028A 1118      BNEQ  70$ ; Nonzero means overflow has occurred
      028D 1119      SOBGTR R0,60$ ; Check for end of this loop
      05 028D 1120      RSB ; This is return path for no overflow
      028E 1121
5B  02 88 028E 1122 70$: BISB  #PSLSM_V,R11 ; Indicate that overflow has occurred
      05 0291 1123      RSB ; ... and return to the caller

```

```

:LJK0045 0292 .1
:LJK0045 0292 .2
:LJK0045 0292 .3
:LJK0045 0292 .4
:LJK0045 0292 .5
:LJK0045 0292 .6
:LJK0045 0292 .7
:LJK0045 0292 .8
:LJK0045 0292 .9
:LJK0045 0292 .10
:LJK0045 0292 .11
:LJK0045 0292 .12
:LJK0045 0292 .13
:LJK0045 0292 .14
:LJK0045 0292 .15
:LJK0045 0292 .16
:LJK0045 0292 .17
:LJK0045 0292 .18
:LJK0045 0292 .19
:LJK0045 0292 .20
:LJK0045 0292 .21
:LJK0045 0292 .22
:LJK0045 0292 .23
:LJK0045 0292 .24
:LJK0045 0292 .25
:LJK0045 0292 .26
:LJK0045 0292 .27
:LJK0045 0292 .28
:LJK0045 0292 .29
:LJK0045 0292 .30
:LJK0045 0292 .31
:LJK0045 0292 .32
:LJK0045 0292 .33
:LJK0045 0292 .34
:LJK0045 0292 .35
:LJK0045 0292 .36
:LJK0045 0292 .37
:LJK0045 0292 .38
:LJK0045 0292 .39
:LJK0045 0292 .40
:LJK0045 0292 .41
:LJK0045 0292 .42
:LJK0045 0292 .43
:LJK0045 0292 .44
:LJK0045 0292 .45
:LJK0045 0292 .46
:LJK0045 0292 .47
:LJK0045 0292 .48
:LJK0045 0292 .49
:LJK0045 0292 .50
:LJK0045 0292 .51
:LJK0045 0292 .52
:LJK0045 0292 .53
:LJK0045 0292 .54
:LJK0045 0292 .55
:LJK0045 0292 .56
:LJK0045 0292 .57

```

.SUBTITLE - CHECK\_WRITE\_ACCESS - Check Writability of Decimal String

Functional Description:

The ADDP4 and SUBP4 instructions are unique in that they are the only two instructions that read and write the same packed decimal string. They are, in fact, implemented as ADDP6 and SUBP6 where the second input string, ADDend2 or MINuend, and the result string, SUM or DIFFerence, are the same. But ADDP6 and SUBP6, as well as all other packed decimal instructions except ADDP4 and SUBP4, produce UNPREDICTABLE results when an output string overlaps any input string. With this interpretation, ADDP4 and SUBP4 are the only packed decimal instructions that permit overlapping packed decimal strings. (Note that the result string, SUM or DIFFerence, may not overlap the first input string, ADDend or SUBtrahend.)

The implementation of ADDP4 and SUBP4, interpreted as ADDP6 and SUBP6 with overlapping strings, needs to protect itself from modifying memory until the entire instruction can execute to completion. Otherwise, the instruction will be restarted with a different initial state than it first had. This is accomplished by probing the result string for write access before execution begins. This routine receives control if that PROBEW fails.

In the interest of simplicity, the PROBEW that is executed always uses 16 as the byte count of the result string. This routine can then be called not only when the output string is inaccessible but also when

The output string is shorter than 16 bytes.

The address of the output string is within 16 bytes of the end of the page.

The page containing the string is writable.

The next page is not writable.

This routine distinguishes inaccessible strings from this rare case of a PROBEW failure.

In other words, this routine checks the write accessibility of packed decimal strings. If the entire string is writable, control is passed back to the caller where emulation continues. If the string is not writable, an access violation is generated.

Input Parameters:

R4 - Digit count of result string  
R5 - Address of most significant digit in result string

Output Parameters:

None

Implicit Output:

If the string with its correct byte count is writable, control

```

:LJK0045      0292      .58      :      is returned to the caller.
:LJK0045      0292      .59      :
:LJK0045      0292      .60      :      If the string with its correct digit count is not writable, then the
:LJK0045      0292      .61      :      string is accessed to force an access violation to occur. This will
:LJK0045      0292      .62      :      cause control to be transferred eventually to the access violation
:LJK0045      0292      .63      :      handler at the end of this module that will back out the instruction
:LJK0045      0292      .64      :      so that it can be restarted.
:LJK0045      0292      .65      :
:LJK0045      0292      .66      : Assumption:
:LJK0045      0292      .67      :
:LJK0045      0292      .68      :      This routine assumes that the check for illegal digit count has
:LJK0045      0292      .69      :      already been made so that R4 is between 0 and 31 inclusive.
:LJK0045      0292      .70      :
:LJK0045      0292      .71      : Note:
:LJK0045      0292      .72      :
:LJK0045      0292      .73      :      It is necessary to actually touch the inaccessible page so that the
:LJK0045      0292      .74      :      correct reason mask is generated. Logic that uses the PROBEW
:LJK0045      0292      .75      :      instruction can determine the inaccessible virtual address but is
:LJK0045      0292      .76      :      unable to distinguish, for example, between length violations and page
:LJK0045      0292      .77      :      protection violations.
:LJK0045      0292      .78      :-
:LJK0045      0292      .79      :
:LJK0045      0292      .80      CHECK_WRITE_ACCESS:
:LJK0045      54      04      01      DL      0292      .81      PUSHL      R4      ; Save the digit count
:LJK0045      54      04      01      EF      0294      .82      EXTZV      #1,#4,R4 ; Convert digits to bytes
:LJK0045      54      04      01      EF      0298      :
:LJK0045      65      54      00      D6      0299      .83      INCL      R4      ; Count the byte that contains the sign
:LJK0045      65      54      00      OD      029B      .84      PROBEW    #0,R4,(R5) ; Check writability of the string
:LJK0045      54      8E      04      13      029F      .85      BEQL      10$      ; Branch if string cannot be written
:LJK0045      54      8E      05      D0      02A1      .86      MOVL      (SP)+,R4 ; Restore saved R4
:LJK0045      54      8E      05      05      02A4      .87      RSB      ; Return to caller. String is OK
:LJK0045      54      8E      05      05      02A5      .88      :
:LJK0045      54      8E      05      05      02A5      .89      : The first and last bytes of the string are touched (written) with an
:LJK0045      54      8E      05      05      02A5      .90      :
:LJK0045      54      8E      05      05      02A5      .91      :      ADDB2      #0,xxx
:LJK0045      54      8E      05      05      02A5      .92      :
:LJK0045      54      8E      05      05      02A5      .93      : instruction. This instruction causes the correct access violation reason
:LJK0045      54      8E      05      05      02A5      .94      : mask to be generated but does not modify the contents of locations that are
:LJK0045      54      8E      05      05      02A5      .95      : accessible. Note that at least one of the following two ADDB2 instructions
:LJK0045      54      8E      05      05      02A5      .96      : is guaranteed to generate an access violation, transferring control in a
:LJK0045      54      8E      05      05      02A5      .97      : rather complicated way to the ADD_SUB_BSBW_4 access violation handler.
:LJK0045      54      8E      05      05      02A5      .98      :
:LJK0045      65      00      80      02A5      .99      MARK POINT      ADD_SUB_BSBW_4
:LJK0045      65      00      80      02A5      .100     10$:      ADDB2      #0,(R5) ; Touch the first byte
:LJK0045      6544     00      80      02A8      .101     MARK POINT      ADD_SUB_BSBW_4
:LJK0045      6544     00      80      02A8      .102     ADDB2      #0,(R5)[R4] ; Touch the last byte
:LJK0045      6544     00      80      02A8      .103     :
:LJK0045      6544     00      80      02AC      .104     : We should never reach here unless the PROBE instruction is broken. We
:LJK0045      6544     00      80      02AC      .105     : will leave this code path in for now but remove it at the same time we
:LJK0045      6544     00      80      02AC      .106     : change the other two HALT instructions in this module into software
:LJK0045      6544     00      80      02AC      .107     : generated machine checks exceptions.
:LJK0045      6544     00      80      02AC      .108     :
:LJK0045      6544     00      80      02AC      .109     :: ***** BEGIN TEMP *****
:LJK0045      6544     00      80      02AC      .110     ::
:LJK0045      6544     00      80      02AC      .111     :: THE FOLLOWING HALT INSTRUCTION SHOULD BE REPLACED WITH THE CORRECT
:LJK0045      6544     00      80      02AC      .112     :: ABORT CODE.
:LJK0045      6544     00      80      02AC      .113     ::

```

VAX\$DECIMAL\_ARITHMETIC  
V04-001

F 9

- VAX-11 Packed Decimal Arithmetic Instr 8-JAN-1985 17:27:01 VAX/VMS Macro V04-00 Page 28  
- CHECK\_WRITE\_ACCESS - Check Writability 5-SEP-1984 00:44:34 [EMULAT.BUGSRC]VAXARITH.MAR;1 (13)

:LJK0045  
:LJK0045  
:LJK0045

00 02AC .114 halt ; This will cause an OPCDEC exception  
02AD .115 :::  
02AD .116 ::: \*\*\*\*\* END TEMP \*\*\*\*\*

```
02AD 1125 .SUBTITLE VAX$MULP - Multiply Packed
02AD 1126
02AD 1127 : Functional Description:
02AD 1128 :
02AD 1129 : The multiplicand string specified by the multiplicand length and
02AD 1130 : multiplicand address operands is multiplied by the multiplier string
02AD 1131 : specified by the multiplier length and multiplier address operands. The
02AD 1132 : product string specified by the product length and product address
02AD 1133 : operands is replaced by the result.
02AD 1134
02AD 1135 : Input Parameters:
02AD 1136 :
02AD 1137 : R0 - mulrlen.rw Number of digits in multiplier string
02AD 1138 : R1 - mulraddr.ab Address of multiplier string
02AD 1139 : R2 - mulrlen.rw Number of digits in multiplicand string
02AD 1140 : R3 - muldaddr.ab Address of multiplicand string
02AD 1141 : R4 - prodlen.rw Number of digits in product string
02AD 1142 : R5 - prodaddr.ab Address of product string
02AD 1143
02AD 1144 : Output Parameters:
02AD 1145 :
02AD 1146 : R0 = 0
02AD 1147 : R1 = Address of the byte containing the most significant digit of
02AD 1148 : the multiplier string
02AD 1149 : R2 = 0
02AD 1150 : R3 = Address of the byte containing the most significant digit of
02AD 1151 : the multiplicand string
02AD 1152 : R4 = 0
02AD 1153 : R5 = Address of the byte containing the most significant digit of
02AD 1154 : the string containing the product
02AD 1155
02AD 1156 : Condition Codes:
02AD 1157 :
02AD 1158 : N <- product string LSS 0
02AD 1159 : Z <- product string EQL 0
02AD 1160 : V <- decimal overflow
02AD 1161 : C <- 0
02AD 1162
02AD 1163 : Register Usage:
02AD 1164 :
02AD 1165 : This routine uses all of the general registers. The condition codes
02AD 1166 : are computed at the end of the instruction as the final result is
02AD 1167 : stored in the product string. R11 is used to record the condition
02AD 1168 : codes.
02AD 1169
02AD 1170 : Notes:
02AD 1171 :
02AD 1172 : 1. This routine uses a large amount of stack space to allow storage of
02AD 1173 : intermediate results in a convenient form. Specifically, each digit
02AD 1174 : pair of the longer input string is stored in binary in a longword on
02AD 1175 : the stack. In addition, 32 longwords are set aside to hold the product
02AD 1176 : intermediate result. Each longword contains a binary number between 0
02AD 1177 : and 99.
02AD 1178
02AD 1179 : After the multiplication is complete, Each longword is removed from
02AD 1180 : the stack, converted to a packed decimal pair, and stored in the
02AD 1181 : output string. Any nonzero cells remaining on the stack after the
```



```

02AD 1182 : output string has been completely filled are the indication of decimal
02AD 1183 : overflow.
02AD 1184 :
02AD 1185 : The purpose of this method of storage is to avoid decimal/binary or
02AD 1186 : even byte/longword conversions during the calculation of intermediate
02AD 1187 : results.
02AD 1188 :
02AD 1189 : 2. Trailing zeros are removed from the larger string. All zeros in
02AD 1190 : the shorter string are eliminated in the sense that no arithmetic
02AD 1191 : is performed. The output array pointer is simply advanced to point
02AD 1192 : to the next higher array element.
02AD 1193 :-
02AD 1194 :-
02AD 1195 VAX$MULP::
OFFF 8F BB 02AD 1196 PUSHR #*M<R0,R1,R2,R3,R4,R5,R6,R7,R8,R9,R10,R11> ; Save the lot
02B1 1197
02B1 1198 ESTABLISH HANDLER - ; Store address of access
02B1 1199 ARITH_ACCVIO ; violation handler
02B6 1200
02B6 1201 ROPRAND_CHECK R4 ; Insure that R4 is LEQU 31
02C1 1202
02C1 1203 ROPRAND_CHECK R2 ; Insure that R2 is LEQU 31
02C9 1204 MARK_POINT MULP BSBW 0
FD34' 30 02C9 1205 BSBW DECIMAL$STRIP_ZEROS_R2_R3 ; Strip high order zeros from R2/R3
02CC 1206
02CC 1207 ROPRAND_CHECK R0 ; Insure that R0 is LEQU 31
02D4 1208 MARK_POINT MULP BSBW 0
FD29' 30 02D4 1209 BSBW DECIMAL$STRIP_ZEROS_R0_R1 ; Strip high order zeros from R0/R1
02D7 1210
50 04 01 EF 02D7 1211 EXTZV #1,#4,R0,R0 ; Convert digit count to byte count
02DB 1212
50 06 D6 02DC 1212 INCL R0 ; Include least significant digit
02DE 1213
52 04 01 EF 02DE 1214 EXTZV #1,#4,R2,R2 ; Convert digit count to byte count
02E2 1215
52 06 D6 02E3 1215 INCL R2 ; Include least significant digit
02E5 1216
52 50 D1 02E5 1217 CMPL R0,R2 ; See which string is larger
02E8 1218 BGTRU 3$ ; R2/R3 describes the longer string
58 52 7D 02EA 1219 MOVQ R2,R8 ; R8 and R9 describe the longer string
7E 50 7D 02ED 1220 MOVQ R0,-(SP) ; Shorter string descriptor also saved
06 11 02F0 1221 BRB 6$
02F2 1222
58 50 7D 02F2 1223 3$: MOVQ R0,R8 ; R8 and R9 describe the longer string
7E 52 7D 02F5 1224 MOVQ R2,-(SP) ; Shorter string descriptor also saved
02F8 1225
02F8 1226 ; Create space for the output array on the stack (32 longwords of zeros)
02F8 1227
50 08 D0 02F8 1228 6$: MOVL #8,R0 ; Eight pairs of quadwords
02FB 1229
7E 7C 02FB 1230 10$: CLRQ -(SP) ; Clear one pair
7E 7C 02FD 1231 CLRQ -(SP) ; ... and another
F9 50 F5 02FF 1232 SOBGTR R0,10$ ; Do all eight pairs
0302 1233
57 5E D0 0302 1234 MOVL SP,R7 ; Store beginning of output array in R7
0305 1235
0305 1236 ; The longer input array will be stored on the stack as an array of

```

```

0305 1237 : longwords. Each array element contains a number between 0 and 99,
0305 1238 : representing a pair of digits in the original packed decimal string.
0305 1239 : Because the units digit is stored with the sign in packed decimal format,
0305 1240 : it is necessary to shift the number as we store it. This is accomplished by
0305 1241 : multiplying the number by ten.
0305 1242 :
0305 1243 : The longer array is described by R8 (byte count) and R9 (address of most
0305 1244 : significant digit pair).
0305 1245 :
55 58 59 C1 0305 1246 ADDL3 R9,R8,R5 ; Point R5 beyond sign digit
54 58 D0 0309 1247 MOVL R8,R4 ; R4 contains the loop count
030C 1248 :
030C 1249 : An array of longwords is allocated on the stack. R3 starts out pointing
030C 1250 : at the longword beyond the top of the stack. The first remainder, guaranteed
030C 1251 : to be zero, is "stored" here. The rest of the digit pairs are stored safely
030C 1252 : below the top of the stack.
030C 1253 :
53 53 58 CE 030C 1254 MNEGL R8,R3 ; Stack grows toward lower addresses
5E 6E43 DE 030F 1255 MOVAL (SP)[R3],SP ; Allocate the space
53 5E 04 C3 0313 1256 SUBL3 #4,SP,R3 ; Point R3 at next lower longword
0317 1257 :
0317 1258 MARK POINT MULP_R8
51 51 75 9A 0317 1259 20$: MOVZBL -(R5),R1 ; Get next digit pair
0000'CF41 9A 031A 1260 MOVZBL DECIMAL$PACKED_TO_BINARY_TABLE[R1],-
0320 1261 R1 ; Convert digits to binary
52 51 0A 7A 0320 1262 EMUL #10,R1,R2,R0 ; Multiply by 10
50 0324 :
00000064 8F 7B 0325 1263 EDIV #100,R0,R2,(R3)+ ; Divide by 100
83 52 50 032B :
E6 54 F5 032E 1264 SOBGTR R4,20$
0331 1265 :
63 52 D0 0331 1266 MOVL R2,(R3) ; Store final quotient
59 5E D0 0334 1267 MOVL SP,R9 ; Remember array address in R9
6E48 DF 0337 1268 PUSHAL (SP)[R8] ; Store start of fixed size area
033A 1269 :
033A 1270 : Check for trailing zeros in the input array stored on the stack. If any are
033A 1271 : present, they are removed and the product array is adjusted accordingly.
033A 1272 :
89 D5 033A 1273 30$: TSTL (R9)+ ; Is next number zero?
08 12 033C 1274 BNEQ 40$ ; Leave loop if nonzero
57 04 C0 033E 1275 ADDL #4,R7 ; Advance output pointer to next element
F6 58 F5 0341 1276 SOBGTR R8,30$ ; Keep going
0344 1277 :
0344 1278 : If we drop through the loop, then the entire input array is zero. There is
0344 1279 : no need to perform any arithmetic because the product will be zero (and the
0344 1280 : output array on the stack starts out as zero). The only remaining work is
0344 1281 : to store the result in the output string and set the condition codes.
0344 1282 :
2C 11 0344 1283 BRB 70$ ; Exit to end processing
0346 1284 :
0346 1285 : Now multiply the input array by each successive digit pair. In order to
0346 1286 : allow R10 to continue to locate ARITH_ACCVIO while we execute this loop, it
0346 1287 : is necessary to perform a small amount of register juggling. In essence,
0346 1288 : R8 and R9 switch the identity of the string that they describe.
0346 1289 :
59 04 C2 0346 1290 40$: SUBL #4,R9 ; Readjust input array pointer
7E 5E 7D 0349 1291 MOVQ R8,-(SP) ; Save R8/R9 descriptor on stack

```

58	08	AE	D0	034C	1292	MOVL	8(SP),R8	:	Point R8 at start of 32-longword array
58	0080	C8	7D	0350	1293	MOVQ	<32*4>(R8),R8	:	Get descriptor that follows that array
	59	58	C0	0355	1294	ADDL2	R8,R9	:	Point R9 beyond sign byte
				0358	1295				
	53	87	DE	0358	1296	50\$:	MOVAL (R7)+,R3	:	Output array address to R3
				035B	1297		MARK POINT MULP_AT_SP		
	51	79	9A	035B	1298		MOVZBL -(R9),R1	:	Next digit pair to R1
56	0000	'CF41	9A	035E	1299		MOVZBL DECIMALSPACKED_TO_BINARY_TABLE[R1],-		
				0364	1300		R6	:	Convert digits to binary
		06	13	0364	1301		BEQL 60\$	:	Skip the work if zero
	54	6E	7D	0366	1302		MOVQ (SP),R4	:	Input array descriptor to R4/R5
		0104	30	0369	1303		BSBW EXTEND_STRING_MULTIPLY	:	Do the work
		E9	58	036C	1304	60\$:	SOBGTR R8,50\$	:	Any more multiplier digits?
				036F	1305				
	5E	08	C0	036F	1306		ADDL #8,SP	:	Discard saved long string descriptor
				0372	1307				
	5E	6E	D0	0372	1308	70\$:	MOVL (SP),SP	:	Remove input array from stack
				0375	1309				
				0375	1310			:	At this point, the product string is located in a 32-longword array on
				0375	1311			:	the top of the stack. Each longword corresponds to a pair of digits in
				0375	1312			:	the output string. As digits are removed from the stack, they are checked
				0375	1313			:	for nonzero to obtain the correct setting of the Z-bit. After the output
				0375	1314			:	string has been filled, the remainder of the product string is removed from
				0375	1315			:	the stack. If a nonzero result is detected at this stage, the V-bit is set.
				0375	1316				
	59	20	D0	0375	1317		MOVL #32,R9	:	Set up array counter
54	0098	CE	7D	0378	1318		MOVQ < <32*4> + -	:	Skip over 32-longword array
				037D	1319		<2*4> + -	:	and saved string descriptor
				037D	1320		<4*4> >(SP),R4	:	to retrieve original R4 and R5

```

037D 1322      .SUBTITLE      Common Exit Path for VAXSMULP and VAXSDIV
037D 1323      :
037D 1324      : The code for VAXSMULP and VAXSDIV merges at this point. The result is stored
037D 1325      : in an array of longwords at the top of the stack. The size of this array is
037D 1326      : stored in R9. The original R4 and R5 have been retrieved from the stack.
037D 1327      :
037D 1328      : Input Parameters:
037D 1329      :
037D 1330      :   R4 - (contains byte count of destination string in R4 <1:4>)
037D 1331      :   R5 - Address of most significant digit of destination string
037D 1332      :   R9 - Count of longwords in result array on stack
037D 1333      :
037D 1334      : Contents of result array
037D 1335      :
037D 1336      : Implicit Input:
037D 1337      :
037D 1338      :   Signs of two input factors (multiplier and multiplicand or
037D 1339      :   divisor and dividend)
037D 1340      : -
037D 1341      :
037D 1342      MULTIPLY DIVIDE_EXIT:
037D 1343      MOVPSL R11          : Get current PSL
04  00  04  DC 037F 1344      INSV #PSLSM_Z,#0,#4,R11      : Clear all codes except Z-bit
0383      :
0384 1345      ESTABLISH HANDLER -          : Store address of access
0384 1346      ARITH_ACCVIO          : violation handler again
54  04  01  EF 0309 1347      EXTZV #1,#4,R4,R3      : Excess byte count to R3
038D      :
038E 1348      BEQL 125$          : Skip to single digit code
57  55  53  C1 0390 1349      ADDL3 R3,R5,R7          : Remember address of sign byte
55  57  01  C1 0394 1350      ADDL3 #1,R7,R5          : Point R5 beyond end of product string
0398 1351      :
0398 1352 80$: MOVL (SP)+,R1          : Remove next value from stack
0398 1353      BEQL 90$          : Do not clear Z-bit if zero
039D 1354      BICB2 #PSLSM_Z,R11      : Clear Z-bit
03A0 1355      :
03A0 1356      MARK_POINT      MULDIVP_R9
75  0000'CF41 90 03A0 1357 90$: MOVB DECIMAL$BINARY_TO_PACKED_TABLE[R1],-
03A6 1358      -(R5)          : Store converted sum byte
03A6 1359      DECL R9          : One less element on the stack
03A8 1360      BLEQ 116$        : Exit loop if result array exhausted
EB  53  F5 03AA 1361      SOBGTR R3,80$      : Keep going?
03AD 1362      :
03AD 1363 100$: BLBC R4,120$      : Different for even digit count
0380 1364      :
0380 1365      : The output string consists of an odd number of digits. A complete digit
0380 1366      : pair can be stored in the most significant (lowest addressed) byte of
0380 1367      : the product string.
0380 1368      :
0380 1369      MOVL (SP)+,R1          : Remove next value from stack
0383 1370      BEQL 110$        : Do not clear Z-bit if zero
0385 1371      BICB2 #PSLSM_Z,R11      : Clear Z-bit
0388 1372      :
0388 1373      MARK_POINT      MULDIVP_R9
75  0000'CF41 90 0388 1374 110$: MOVB DECIMAL$BINARY_TO_PACKED_TABLE[R1],-
038E 1375      -(R5)          : Store converted sum byte
038E 1376      DECL R9          : One less element on the stack

```

```

04 15 03C0 1377 BLEQ 1168 ; Exit loop if result array exhausted
38 11 03C2 1378 BRB 1408 ; Perform overflow check
03C4 1379
03C4 1380 ; This loop executes if the result array has fewer elements than the output
03C4 1381 ; string. The remaining bytes in the output string are filled with zeros.
03C4 1382 ; There is no need for an overflow check.
03C4 1383
03C4 1384 MARK POINT MULP_DIVP_8
75 94 03C4 1385 114$: CLRB -(R5) ; Store another zero byte
FB 53 F4 03C6 1386 116$: SOBGEQ R3,114$ ; Any more room in output string
03C9 1387
38 11 03C9 1388 BRB 150$ ; Determine sign of result
03CB 1389
03CB 1390 ; This code path is used in the case where the output digit count is 0 or 1.
03CB 1391 ; R5 must be advanced
03CB 1392
57 55 D0 03CB 1393 125$: MOVL R5,R7 ; Remember address of output sign byte
55 D6 03CE 1394 INCL R5 ; Advance R5 so common code can be used
DB 11 03D0 1395 BRB 100$ ; Join common code path
03D2 1396
03D2 1397 ; The output string consists of an even number of digits. Only the low order
03D2 1398 ; nibble is stored in the most significant (lowest addresses) byte. A zero is
03D2 1399 ; stored in the high order nibble. If the high order digit would have been
03D2 1400 ; nonzero, the V-bit is set and the overflow check is bypassed because there
03D2 1401 ; are faster ways to clean the stack if we do not have to check for nonzero
03D2 1402 ; at the same time.
03D2 1403
51 51 8E D0 03D2 1404 120$: MOVL (SP)+,R1 ; Remove next value from stack
0000'CF41 90 03D5 1405 MOVB DECIMALSBINARY_TO_PACKED_TABLE[R1],-
03DB 1406 R1 ; Obtain converted sum byte
03DB 1407 MARK POINT MULP_DIVP_R9
51 F0 8F 88 03DB 1408 BICB3 #^XF0,R1,-(R5) ; Store byte, clearing high order nibble
03DF 1409
03 13 03E0 1409 BEQL 130$ ; Do not clear Z-bit if zero
51 5B 04 8A 03E2 1410 BICB2 #PSLSM Z,R11 ; Clear Z-bit
F0 8F 93 03E5 1411 130$: BITB #^XF0,R1 ; Is high order nibble nonzero?
06 12 03E9 1412 BNEQ 133$ ; Yes, go set overflow bit
59 D7 03EB 1413 DECL R9 ; One less element on the stack
D7 15 03ED 1414 BLEQ 116$ ; Exit loop if result array exhausted
0B 11 03EF 1415 BRB 140$ ; Check rest of result array for nonzero
03F1 1416
03F1 1417 ; If we detect overflow, we need to adjust R9 to reflect the nonzero longword
03F1 1418 ; removed from the stack before we enter the next code block that sets the
03F1 1419 ; V-bit and cleans off the stack based on the contents of R9.
03F1 1420
59 D7 03F1 1421 133$: DECL R9 ; One more longword removed from stack
03F3 1422
03F3 1423 ; A nonzero digit has been discovered in a position that cannot be stored in
03F3 1424 ; the output string. Set the V-bit, remove the rest of the product array from
03F3 1425 ; the stack, and join the exit processing in the code that determines the sign
03F3 1426 ; of the product.
03F3 1427
SE 5B 02 88 03F3 1428 135$: BISB #PSLSM V,R11 ; Set the overflow bit
6E49 DE 03F6 1429 MOVAL (SP)[R9],SP ; Clean off remaining product string
07 11 03FA 1430 BRB 150$ ; Go to code that determines the sign
03FC 1431
03FC 1432 ; The remainder of the product array must be removed from the stack. A nonzero

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```

03FC 1433 ; result causes the V-bit to be set and the rest of the loop to be skipped.
03FC 1434 ; Note that there is always a nonzero loop count remaining at this point.
03FC 1435
      BE D5 03FC 1436 140$: TSTL (SP)+ ; Is next longword zero?
      F1 12 03FE 1437 BNEQ 133$ ; No, leave loop
      F9 59 F5 0400 1438 SOBGR R9,140$
      0403 1439
      0403 1440 ; The final product string has been stored and the V- and Z-bits have their
      0403 1441 ; correct settings. The sign of the product must be determined from the
      0403 1442 ; signs of the two input strings. Opposite signs produce a negative product.
      0403 1443 ; Same signs (in any representation) produce a plus sign in the output string.
      0403 1444
      SE 08 C0 0403 1445 150$: ADDL #8,SP ; Discard saved string descriptor
      S6 0C D0 0406 1446 MOVL #12,R6 ; Assume final result is positive
      50 50 6E 7D 0409 1447 MOVQ (SP),R0 ; Retrieve original R0/R1 pair
      04 01 EF 040C 1448 EXTZV #1,#4,R0,R0 ; Get byte count for first input string
      51 50 C0 0410 1449
      0411 1449 ADDL R0,R1 ; Point R1 to byte containing sign
      61 F0 8F 8B 0414 1450 MARK POINT MULP_DIVP_0
      50 0414 1451 BICB3 #*B11110000,(R1),R0 ; R0 contains the sign "digit"
      0418
      0419 1452
      0419 1453 CASE R0,TYPE=B,LIMIT=#10,<- ; Dispatch on sign digit
      0419 1454 220$,- ; 10 => sign is '+'
      0419 1455 210$,- ; 11 => sign is '-'
      0419 1456 220$,- ; 12 => sign is '+'
      0419 1457 210$,- ; 13 => sign is '-'
      0419 1458 220$,- ; 14 => sign is '+'
      0419 1459 220$,- ; 15 => sign is '+'
      0419 1460 >
      54 01 D0 0429 1461
      02 11 042C 1462 210$: MOVL #1,R4 ; Count a minus sign
      042E 1463 BRB 230$ ; Now check second input sign
      54 D4 042E 1464 220$: CLRL R4 ; No real minus signs so far
      0430 1465
      52 08 AE 7D 0430 1467 230$: MOVQ 8(SP),R2 ; Retrieve original R2/R3 pair
      52 04 01 EF 0434 1468 EXTZV #1,#4,R2,R2 ; Get byte count for second input string
      52 52
      53 52 C0 0439 1469 ADDL R2,R3 ; Point R3 to byte containing sign
      043C 1470 MARK POINT MULP_DIVP_0
      63 F0 8F 8B 043C 1471 BICB3 #*B11110000,(R3),R2 ; R2 contains the sign "digit"
      0440
      0441 1472
      0441 1473 CASE R2,TYPE=B,LIMIT=#10,<- ; Dispatch on sign digit
      0441 1474 250$,- ; 10 => sign is '+'
      0441 1475 240$,- ; 11 => sign is '-'
      0441 1476 250$,- ; 12 => sign is '+'
      0441 1477 240$,- ; 13 => sign is '-'
      0441 1478 250$,- ; 14 => sign is '+'
      0441 1479 250$,- ; 15 => sign is '+'
      0441 1480 >
      0451 1481
      10 54 D6 0451 1482 240$: INCL R4 ; Remember that sign was minus
      09 54 E9 0453 1483 250$: BLBC R4,260$ ; Even parity indicates positive result
      5B 02 E0 0456 1484 BBS #PSLSV_Z,R11,270$ ; Step out of line for minus zero check
      5B 08 88 045A 1485 BISB #PSLSM_N,R11 ; Set N-bit in saved PSW

```

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      56 D6 045D 1486 255$: INCL R6 ; Change sign to minus
      045F 1487
      045F 1488 MARK_POINT MULP_DIVP_0
      04 00 56 F0 045F 1489 260$: INSV R6,#0,#4,(R7) ; Store sign in result string
      67 0463
      10 AE D4 0464 1490 CLRL 16(SP) ; Set saved R4 to zero
      FB96 31 0467 1491 BRW VAX$DECIMAL_EXIT ; Join common exit code
      046A 1492
      046A 1493 ; If the result is negative zero, then it must be changed to positive zero
      046A 1494 ; unless overflow has occurred, in which case, the sign is left as negative
      046A 1495 ; but the N-bit is clear.
      EF 5B 01 E0 046A 1497 270$: BBS #PSLSV_V,R11,255$ ; Make sign negative if overflow
      EF 11 046E 1498 BRB 260$ ; Sign will be positive
  
```

```

0470 1500 .SUBTITLE EXTEND_STRING_MULTIPLY - Multiply a String by a Number
0470 1501 :
0470 1502 : Functional Description:
0470 1503 :
0470 1504 : This routine multiplies an array of numbers (each array element LEQU
0470 1505 : 99) by a number (also LEQU 99). The resulting product array is added
0470 1506 : to another array, each of whose elements is also LEQU 99.
0470 1507 :
0470 1508 : Input Parameters:
0470 1509 :
0470 1510 : R3 - Pointer to output array
0470 1511 : R4 - Input array size
0470 1512 : R5 - Input array address
0470 1513 : R6 - Multiplier
0470 1514 :
0470 1515 : Output Parameters:
0470 1516 :
0470 1517 : None
0470 1518 :
0470 1519 : Implicit Output:
0470 1520 :
0470 1521 : The output array is altered.
0470 1522 :
0470 1523 : An intermediate product array is produced by multiplying each input
0470 1524 : array element by the multiplier. Each product array element is then
0470 1525 : added to the corresponding output array element.
0470 1526 :
0470 1527 : Side Effects:
0470 1528 :
0470 1529 : R3, R4, and R5 are modified by this routine.
0470 1530 :
0470 1531 : R6 is preserved.
0470 1532 :
0470 1533 : R0, R1, and R2 are used as scratch registers. R0 and R1 contain the
0470 1534 : quadword result of EMUL that is then passed into EDIV.
0470 1535 :
0470 1536 : Assumptions:
0470 1537 :
0470 1538 : This routine assumes that all array elements lie in the range from 0
0470 1539 : to 99 inclusive. (This is true if all input strings contain only legal
0470 1540 : decimal digits.) The arithmetic performed by this routine will
0470 1541 : maintain this assumption. That is,
0470 1542 :
0470 1543 :
0470 1544 : times      input array element      LEQU 99
0470 1545 :            multiplier                LEQU 99
0470 1546 :            -----
0470 1547 : plus      product                    LEQU 99   LEQU 99*99
0470 1548 :            carry
0470 1549 :            -----
0470 1550 : plus      modified product          LEQU 99   LEQU 99*100
0470 1551 :            old output array element
0470 1552 :            -----
0470 1553 :            new output array element          LEQU 99*101 = 9999
0470 1554 :
0470 1555 : A number LEQU 9999, when divided by 100, is guaranteed to produce both
0470 1556 : a quotient and a remainder LEQU 99.
  
```



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0470 1557
0470 1558 EXTEND_STRING_MULTIPLY:
52   D4 0470 1559      CLRL   R2           ; Initial carry is zero
      0472 1560
52   85 55 7A 0472 1561 10$:  EMUL   R6,(R5)+,R2,R0 ; Form modified product (R0 LEQU 9900)
      50 50 0476
      50 63 C0 0477 1562      ADDL2  (R3),R0      ; Add old output array element
00000064 8F 7B 047A 1563      EDIV   #100,R0,R2,(R3)+ ; Remainder to output array
83   52 50
      EC 54 F5 0480 1564
      0483 1565      SOBGRP  R4,10$      ; Quotient becomes carry
      0486 1566      ; Keep going?
      0486 1567 ; This remaining code looks more complicated than it actually is. In the
      0486 1568 ; usual case, the routine exits immediately. In the event that a carry
      0486 1569 ; occurs, one additional entry in the output array will be modified. Only in
      0486 1570 ; the rare case of an output array consisting of a string of 99s will any
      0486 1571 ; significant looping occur.
      0486 1572
      63 52 C0 0486 1573      ADDL2  R2,(R3)      ; Add final carry
00000064 63 D1 0489 1574 20$:  CMPL   (R3),#100    ; Do we overflow into next digit pair?
      8F 048B
      01 1E 0490 1575      BGEQU  30$      ; Branch if carry
      05 0492 1576      RSB           ; Otherwise, all done
00000064 8F C2 0493 1577
      83 0493 1578 30$:  SUBL   #100,(R3)+ ; Readjust entry and advance pointer
      63 D6 049A 1579      INCL   (R3)      ; Propagate carry
      EB 11 049C 1580      BRB    20$      ; ... and test this entry for overflow
  
```

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049E 1582 .SUBTITLE VAX$DIVP - Divide Packed
049E 1583
049E 1584 : Functional Description:
049E 1585 :
049E 1586 : The dividend string specified by the dividend length and dividend
049E 1587 : address operands is divided by the divisor string specified by the
049E 1588 : divisor length and divisor address operands. The quotient string
049E 1589 : specified by the quotient length and quotient address operands is
049E 1590 : replaced by the result.
049E 1591
049E 1592 : Input Parameters:
049E 1593 :
049E 1594 : R0 - divrlen.rw Number of digits in divisor string
049E 1595 : R1 - divraddr.ab Address of divisor string
049E 1596 : R2 - divdlen.rw Number of digits in dividend string
049E 1597 : R3 - divdaddr.ab Address of dividend string
049E 1598 : R4 - quolen.rw Number of digits in quotient string
049E 1599 : R5 - quoaddr.ab Address of quotient string
049E 1600
049E 1601 : Output Parameters:
049E 1602 :
049E 1603 : R0 = 0
049E 1604 : R1 = Address of the byte containing the most significant digit of
049E 1605 : the divisor string
049E 1606 : R2 = 0
049E 1607 : R3 = Address of the byte containing the most significant digit of
049E 1608 : the dividend string
049E 1609 : R4 = 0
049E 1610 : R5 = Address of the byte containing the most significant digit of
049E 1611 : the string containing the quotient
049E 1612
049E 1613 : Condition Codes:
049E 1614 :
049E 1615 : N <- quotient string LSS 0
049E 1616 : Z <- quotient string EQL 0
049E 1617 : V <- decimal overflow
049E 1618 : C <- 0
049E 1619
049E 1620 : Register Usage:
049E 1621 :
049E 1622 : This routine uses all of the general registers. The condition codes
049E 1623 : are computed at the end of the instruction as the final result is
049E 1624 : stored in the quotient string. R11 is used to record the condition
049E 1625 : codes.
049E 1626
049E 1627 : Algorithm:
049E 1628 :
049E 1629 : This algorithm is the straightforward approach described in
049E 1630 :
049E 1631 : The Art of Computer Programming
049E 1632 : Second Edition
049E 1633 :
049E 1634 : Volume 2 / Seminumerical Algorithms
049E 1635 : Donald E. Knuth
049E 1636 :
049E 1637 : 1981
049E 1638 : Addison-Wesley Publishing Company
  
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Reading, Massachusetts

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049E 1639 :
049E 1640 :
049E 1641 : Notes:
049E 1642 :
049E 1643 : The choice of a longword array to store the auotient deserves a
049E 1644 : comment. In VAX$MULP, a longword array was used because its elements
049E 1645 : were used directly by MULP and DIVP instructions. The use of longwords
049E 1646 : eliminated the need to convert back and forth between longwords and
049E 1647 : bytes. In this routine, the QUOTIENT DIGIT routine returns its result
049E 1648 : in a register, which result can easily be stored in whatever way is
049E 1649 : convenient. By using longwords instead of bytes, this routine can use
049E 1650 : the same end processing code as MULP, a sizeable savings in code.
049E 1651 :-
049E 1652 :
049E 1653 : .ENABLE LOCAL_BLOCK
049E 1654 :
049E 1655 :+
049E 1656 : This code path is entered if the divisor is zero.
049E 1657 :
049E 1658 : Input Parameter:
049E 1659 :
049E 1660 : (SP) - Return PC
049E 1661 :
049E 1662 : Output Parameters:
049E 1663 :
049E 1664 : 0(SP) - SRMSK_FLT_DIV_T (Arithmetic trap code)
049E 1665 : 4(SP) - Final state PSL
049E 1666 : 8(SP) - Return PC
049E 1667 :
049E 1668 : Implicit Output:
049E 1669 :
049E 1670 : Control passes through this code to VAX$REFLECT_TRAP.
049E 1671 :-
049E 1672 :
049E 1673 DIVIDE_BY_ZERO:
OFFF 8F BA 049E 1674 POPR #*M<R0,R1,R2,R3,R4,R5,R6,R7,R8,R9,R10,R11>
04A2 1675 ; Restore registers and reset SP
04A2 1676 MOVPSL -(SP) ; Save final PSL on stack
04 DD 04A4 1677 PUSHL #SRMSK_FLT_DIV_T ; Store arithmetic trap code
FB57 31 04A6 1678 BRW VAX$REFLECT_TRAP ; Report exception
04A9 1679
04A9 1680 ; If the divisor contains more nonzero digits than the dividend, then the
04A9 1681 ; quotient will be identically zero. Set up the stack and the registers (R4,
04A9 1682 ; R5, and R9) so that the exit code will be entered to produce this result.
04A9 1683
04A9 1684 1$: CLRL -(SP) ; Fake a quotient digit
59 7E D4 04AB 1685 MOVL #1,R9 ; Count that digit
01 DO 04AE 1686 BRW MULTIPLY_DIVIDE_EXIT ; Store the zero in the output string
FECC 31 04B1 1687
04B1 1688 VAX$DIVP::
OFFF 8F BB 04B1 1689 PUSHR #*M<R0,R1,R2,R3,R4,R5,R6,R7,R8,R9,R10,R11> ; Save the lot
04B5 1690
04B5 1691 ESTABLISH_HANDLER - ; Store address of access
04B5 1692 ARITH_ACCVIO ; violation handler
04BA 1693
04BA 1694 ROPRAND_CHECK R4 ; Insure that R4 is LEQU 31
04C5 1695

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04C5 1696      ROPRND CHECK   R2      ; Insure that R2 is LEQU 31
04CD 1697      MARK_POINT   DIVP BSBW 0
FB30' 30 04CD 1698      BSBW-  DECIMAL$STRIP_ZEROS_R2_R3 ; Strip high order zeros from R2/R3
04D0 1699
04D0 1700      ROPRND CHECK   R0      ; Insure that R0 is LEQU 31
04D8 1701      MARK_POINT   DIVP BSBW 0
FB25' 30 04D8 1702      BSBW-  DECIMAL$STRIP_ZEROS_R0_R1 ; Strip high order zeros from R0/R1
04DB 1703
04DB 1704      ; Insure that the divisor is not zero. Because leading zeros have already
04DB 1705      ; been eliminated, the divisor can only be zero if R0 is 0 (zero length
04DB 1706      ; strings are identically zero) or 1 (R1 contains a sign digit in the low
04DB 1707      ; order nibble and zero in the high order nibble). Note that an exception
04DB 1708      ; will not be generated if an even length string has an illegal nonzero digit
04DB 1709      ; stored in its most significant nibble (including an illegal form of a zero
04DB 1710      ; length string.
04DB 1711
50 04 01 EF 04DB 1712      EXTZV   #1,#4,R0,R0      ; Convert divisor digit count to bytes
04DF
04E0 1713      BNEQ    10$      ; Skip zero divisor check unless zero
04E2 1714      MARK_POINT   DIVP 0
61  F0 8F 93 04E2 1715      BITB-  #^B11110000,(R1) ; Check for zero in ones digit
04E6 1716      BEQL    DIVIDE_BY_ZERO ; Generate exception if zero
04E8 1717
04E8 1718      ; This routine chooses to do its work with a fair amount of internal storage,
04E8 1719      ; all of it allocated on the stack. The quotient is stored as it is computed,
04E8 1720      ; in a 16-longword array. The dividend and divisor are stored as longword arrays,
04E8 1721      ; with each array element storing a digit pair from the original packed
04E8 1722      ; decimal string. The numerator digits are shifted by one digit (multiplied
04E8 1723      ; by ten) so that the quotient has its digits correctly placed, leaving room
04E8 1724      ; for a sign in the low order nibble of the least significant byte. A scratch
04E8 1725      ; array is also allocated on the stack to accommodate intermediate results
04E8 1726      ; of the QUOTIENT_DIGIT routine.
04E8 1727
04E8 1728 10$: INCL    R0      ; Include least significant digit
58 50 7D 04EA 1729      MOVQ   R0,R8      ; Let R8 and R9 describe the divisor
04ED 1730
52 04 01 EF 04ED 1731      EXTZV   #1,#4,R2,R2      ; Convert dividend digit count to bytes
04F1
04F2 1732      INCL    R2      ; Include least significant digit
7E 52 7D 04F4 1733      MOVQ   R2,-(SP)    ; Save dividend descriptor on stack
04F7 1734
56 52 50 C3 04F7 1735      SUBL3  R0,R2,R6    ; Calculate main loop count
AC 1F 04FB 1736      BLSSU  1$      ; Quotient will be zero
56 D6 04FD 1737      INCL    R6      ; One extra digit is always there
04FF 1738
04FF 1739      ; Allocate R6 longwords of zero on the stack
04FF 1740
50 56 D0 C4FF 1741      MOVL   R6,R0      ; Let R0 be the loop counter
7E D4 0502 1742 15$: CLRL  -(SP)    ; Set aside another quotient digit
FB 50 F5 0504 1743      SOBGTR R0,15$    ; Keep going
0507 1744
57 5E D0 0507 1745      MOVL   SP,R7      ; Remember where this array starts
050A 1746
050A 1747      ; The divisor will be stored on the stack as an array of
050A 1748      ; longwords. Each array element contains a number between 0 and 99,
050A 1749      ; representing a pair of digits in the original packed decimal string.
050A 1750      ; Because the units digit is stored with the sign in packed decimal format.

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050A 1751 ; it is necessary to shift the number as we store it. This is accomplished by
050A 1752 ; multiplying the number by ten.
050A 1753 ;
050A 1754 ; The divisor string is described by R8 (byte count) and R9 (address of most
050A 1755 ; significant digit pair).
050A 1756 ;
55 58 59 C1 050A 1757 ADDL3 R9,R8,R5 ; Point R5 beyond sign digit
54 58 D0 050E 1758 MOVL R8,R4 ; R4 contains the loop count
0511 1759 ;
0511 1760 ; Put in an extra digit place for the divisor. This allows several common
0511 1761 ; subroutines to be used when operating on the divisor string.
0511 1762 ;
7E D4 0511 1763 CLRL -(SP) ; Set aside a place holder
0513 1764 ;
0513 1765 ; An array of longwords is allocated on the stack. R3 starts out pointing
0513 1766 ; at the longword beyond the top of the stack. The first remainder, guaranteed
0513 1767 ; to be zero, is "stored" here. The rest of the digit pairs are stored safely
0513 1768 ; below the top of the stack.
0513 1769 ;
53 53 58 CE 0513 1770 MNEGL R8,R3 ; Stack grows toward lower addresses
SE SE 6E43 DE 0516 1771 MOVAL (SP)[R3],SP ; Allocate the space
53 SE 04 C3 051A 1772 SUBL3 #4,SP,R3 ; Point R3 at next lower longword
051E 1773 ;
051E 1774 ;
51 51 75 9A 051E 1775 20$: MARK POINT DIVP_R6_R7
0000'CF41 9A 0521 1776 MOVZBL -(R5),R1 ; Get next digit pair
0527 1777 MOVZBL DECIMAL$PACKED_TO_BINARY_TABLE[R1],- ; Convert digits to binary
52 51 0A 7A 0527 1778 EMUL #10,R1,R2,R0 ; Multiply by 10
50 052B ;
00000064 8F 7B 052C 1779 EDIV #100,R0,R2,(R3)+ ; Divide by 100
83 52 50 0532 ;
E6 54 F5 0535 1780 SOBGTR R4,20$
0538 1781 ;
0538 1782 ; There are two cases where the final quotient (contents of R2) is zero.
0538 1783 ; In these cases, the number of nonzero digit pairs in the divisor array is
0538 1784 ; smaller by one than the number of bytes containing the original packed decimal
0538 1785 ; string. One case is a divisor string with an even number of digits. The
0538 1786 ; second case is a divisor string with an odd number of digits but the most
0538 1787 ; significant digit is zero (essentially a variant of the first case). The
0538 1788 ; simplest way to handle all of these cases is to decrement R8, the divisor
0538 1789 ; counter, if R2 is zero. Note that previous checks for a zero divisor
0538 1790 ; prevent R8 from going to zero.
0538 1791 ;
63 52 D0 0538 1792 MOVL R2,(R3) ; Store final quotient
0A 12 053B 1793 BNEQ 25$ ; Leave well enough alone if nonzero
56 D6 053D 1794 INCL R6 ; One more quotient digit
57 04 C2 053F 1795 SUBL #4,R7 ; Make room for it
58 D7 0542 1796 DECL R8 ; Count one less divisor "digit"
01 12 0544 1797 BNEQ 25$
0546 1798 ;
0546 1799 :::: ***** BEGIN TEMP *****
0546 1800 ::::
0546 1801 :::: THE FOLLOWING HALT INSTRUCTION SHOULD BE REPLACED WITH THE CORRECT
0546 1802 :::: ABORT CODE.
0546 1803 ::::
0546 1804 :::: THE HALT IS SIMILAR TO THE
0546 1805 ::::

```

```

0546 1806 ::: MICROCODE CANNOT GET HERE
0546 1807 :::
0546 1808 ::: ERRORS THAT OTHER IMPLEMENTATIONS USE.
0546 1809 :::
00 0546 1810 halt ; This will cause an OPCDEC exception
0547 1811 :::
0547 1812 ::: ***** END TEMP *****
0547 1813 :::
59 SE D0 0547 1814 25$: MOVL SP,R9 ; R9 locates low order divisor digit
054A 1815
054A 1816 ; The dividend is stored on the stack as an array of longwords. It does not
054A 1817 ; have its digit pairs shifted so that this storage loop is simpler. An extra
054A 1818 ; place is set aside in the event that it is necessary to normalize the
054A 1819 ; dividend and divisor before division is attempted.
054A 1820
7E D4 054A 1821 CLRL -(SP) ; Set aside space for U[0]
52 6746 DE 054C 1822 MOVAL (R7)[R6],R2 ; Retrieve dividend descriptor
52 62 7D 0550 1823 MOVQ (R2),R2 ; ... in two steps
0553 1824
0553 1825 MARK POINT DIVP_R6_R7
7E 51 83 9A 0553 1826 30$: MOVZBL (R3)+,R1 ; Get next decimal digit pair
0000'CF41 9A 0556 1827 MOVZBL DECIMAL$PACKED_TO_BINARY_TABLE[R1],-
F4 52 F5 055C 1828 -(SP) ; Convert digits to binary
055C 1829 SOBGTR R2,30$ ; Loop through entire input string
055F 1830
055F 1831 ; From this point until the common exit path for MULP and DIVP is entered,
055F 1832 ; no access violations that need to be backed out can occur. We do not need
055F 1833 ; to keep the address of ARITH_ACCVIO in R10 for this stretch of code. Note
055F 1834 ; that R10 must be reloaded before the exit code executes because the
055F 1835 ; destination string is written and may cause access violations.
055F 1836
SA 6746 D0 055F 1837 MOVL (R7)[R6],R10 ; Retrieve size of dividend array
5B SE D0 0563 1838 MOVL SP,R11 ; R11 locates low order dividend digit
0566 1839
0566 1840 ; Allocate a scratch array on the stack the same size as the divisor array
0566 1841 ; (which is one larger than the number of digit pairs)
0566 1842
SE 52 58 CE 0566 1843 MNEGL R8,R2 ; Need a negative index
FC AE42 DE 0569 1844 MOVAL -4(SP)[R2],SP ; Adjust stack pointer
056E 1845
056E 1846 :+
056E 1847 ; At this point, the stack and relevant general registers contain the
056E 1848 ; following information. In this description, N represents the number
056E 1849 ; of digit pairs in the divisor and M represents the number of digit
056E 1850 ; pairs in the dividend.
056E 1851 :
056E 1852 :
056E 1853 : scratch [-----] <-- SP
056E 1854 : [-----] <-- R11
056E 1855 : dividend [-----]
056E 1856 : [-----] <-- R9
056E 1857 : divisor [-----]
056E 1858 : [-----] <-- R7
056E 1859 : quotient [-----]
056E 1860 : [-----]
056E 1861 : [-----]
056E 1862 : [-----]

```

```

056E 1863 :
056E 1864 : R6 - Number of longwords in quotient array (M+1-N)
056E 1865 : R7 - Address of beginning of quotient array
056E 1866 : R8 - Number of digit pairs in divisor (called N)
056E 1867 : R9 - Address of low order digits in divisor
056E 1868 : R10 - Number of digit pairs in dividend (called M)
056E 1869 : R11 - Address of low order digits in dividend
056E 1870 :-
056E 1871 :-
056E 1872 : PUSHAL (SP) ; Store address of scratch array
7E 6E DF 0570 1873 : MOVQ R8,-(SP) ; Remember divisor descriptor
7E 58 7D 0573 1874 : MOVQ R10,-(SP) ; Remember dividend descriptor
7E 5A 7D 0576 1875 :
0576 1876 : The algorithm that guesses the quotient digit can be guaranteed to be off
0576 1877 : by no more than two if the high order digit of the divisor (called V[1]) is
0576 1878 : at least as large as 50 (our radix divided by 2). If the high order digit
0576 1879 : is too small, we 'normalize' the numerator and denominator by multiplying
0576 1880 : them by the same number, namely 100/(V[1]+1).
0576 1881 :
FC A948 01 C1 0576 1882 : ADDL3 #1,-4(R9)[R8],R0 ; Compute V[1] + 1
0576 1883 :
33 50 D1 057C 1883 : CMPL R0,#51 ; Compare to 50 + 1
0576 1884 : BGEQ 40$ ; Skip normalization if V[1] big enough
0576 1885 : DIVL3 R0,#100,R3 ; Compute normalization factor
00000064 8F 0583 :
0583 1886 :
54 58 7D 0589 1886 : MOVQ R8,R4 ; Get descriptor of divisor
00E0 30 058C 1887 : BSBW MULTIPLY_STRING ; Normalize divisor
54 5A 7D 058F 1888 : MOVQ R10,R4 ; Get descriptor of dividend
00DA 30 0592 1889 : BSBW MULTIPLY_STRING ; Normalize dividend
0595 1890 :
0595 1891 : We have now reached the point where we can start calculating quotient digits.
0595 1892 : In the following loop, R5 and R6 are loop invariants. R5 contains the number
0595 1893 : of digit pairs in the divisor. R6 always points to the longword beyond the
0595 1894 : most significant digit in the dividend string. R7 and R8 must be loaded on
0595 1895 : each pass through because these two pointers are modified. Notice that the
0595 1896 : address of the divisor array is exactly what we want to store in R6.
0595 1897 :
5A 56 7D 0595 1898 40$: MOVQ R6,R10 ; Let R10/R11 describe quotient and loop
5A 58 DD 0598 1899 : PUSHL R11 ; Save quotient address for exit code
5B 6B4A DE 059A 1900 : MOVAL (R11)[R10],R11 ; Store quotient digits from high end
059E 1901 :
059E 1902 : This rather harmless looking loop is where the work is done
059E 1903 :
55 58 7D 059E 1904 : MOVQ R8,R5 ; Initialize count and dividend address
59 5A DO 05A1 1905 : MOVL R10,R9 ; Remember the loop count in R9
05A4 1906 :
57 10 AE 7D 05A4 1907 50$: MOVQ 16(SP),R7 ; Load divisor and scratch addresses
001F 30 05AB 1908 : BSBW QUOTIENT_DIGIT ; Get the next quotient digit
7B 53 DO 05AB 1909 : MOVL R3,-(R11) ; Store it
56 04 C2 05AE 1910 : SUBL #4,R6 ; "Advance" dividend pointer
F0 5A F5 05B1 1911 : SOBGTR R10,50$ ; ... and go back for more
05B4 1912 :
05B4 1913 : The quotient digits have been stored on the stack. Eliminate the rest of the
05B4 1914 : stack storage and enter the completion code that this routine shares with
05B4 1915 : VAX$MULP. Note that R9 is already set up with the longword count use by
05B4 1916 : the exit code. Note also that R11 is pointing to the saved dividend descriptor

```





OSCA 1938 .SUBTITLE QUOTIENT\_DIGIT - Get Next Digit in Quotient  
OSCA 1939  
OSCA 1940 : Functional Description:  
OSCA 1941 :  
OSCA 1942 : This routine divides an (N+1)-element array of longwords by an N-element  
OSCA 1943 : array, producing a single quotient digit in the range of 0 to 99  
OSCA 1944 : inclusive. The dividend array is modified by subtracting the product  
OSCA 1945 : of the divisor array and the quotient digit.  
OSCA 1946 :  
OSCA 1947 : The "numbers" that this array operates on multiple precision numbers  
OSCA 1948 : in radix 100. Each digit (a number between 0 and 99) is stored in a  
OSCA 1949 : longword array element with more significant digits stored at higher  
OSCA 1950 : addresses. The dividend string and the scratch string (also called the  
OSCA 1951 : product string) contain one more element than the divisor string.  
OSCA 1952 :  
OSCA 1953 : Input Parameters:  
OSCA 1954 :  
OSCA 1955 : R5 - Number of "digits" (array elements) in divisor array (preserved)  
OSCA 1956 : R6 - Address of longword immediately following most significant  
OSCA 1957 : digit of dividend string (preserved)  
OSCA 1958 : R7 - Address of least significant digit in divisor string (modified)  
OSCA 1959 : R8 - Address of least significant digit in product string (modified)  
OSCA 1960 :  
OSCA 1961 : Output Parameters:  
OSCA 1962 :  
OSCA 1963 : R3 - The quotient that results from dividing the dividend string  
OSCA 1964 : by the divisor string.  
OSCA 1965 :  
OSCA 1966 : The final states of the three pointer registers are listed here  
OSCA 1967 : for completeness.  
OSCA 1968 :  
OSCA 1969 : R6 - Address of longword immediately following most significant  
OSCA 1970 : digit of dividend string  
OSCA 1971 :  
OSCA 1972 : R7 - Address of longword immediately following most significant digit  
OSCA 1973 : of divisor string. This longword must always contain zero.  
OSCA 1974 :  
OSCA 1975 : R8 - Address of longword immediately following most significant  
OSCA 1976 : digit of product string  
OSCA 1977 :  
OSCA 1978 : Implicit Output:  
OSCA 1979 :  
OSCA 1980 : The contents of the dividend array are modified to reflect the  
OSCA 1981 : subtraction of the product string. The result of this subtraction  
OSCA 1982 : could be stored elsewhere. It is a convenience to store it in the  
OSCA 1983 : dividend array on top of those array elements that are no longer  
OSCA 1984 : needed.  
OSCA 1985 :  
OSCA 1986 : The contents of the divisor array are preserved.  
OSCA 1987 :  
OSCA 1988 : Side Effects:  
OSCA 1989 :  
OSCA 1990 : R7 and R8 are modified by this routine. (See implicit output list.)  
OSCA 1991 :  
OSCA 1992 : R5 and R6 are preserved.  
OSCA 1993 :  
OSCA 1994 : R0, R1, R2, and R4 are used as scratch registers. R0 and R1 contain the

```

05CA 1995 : quadword result of EMUL that is then passed into EDIV R2 is the
05CA 1996 : carry from one step to the next. R4 is the loop counter.
05CA 1997 :-
05CA 1998 :-
05CA 1999 QUOTIENT DIGIT:
05CA 2000 EMUL #100,-4(R6),-8(R6),R0 ; R0 <- 100 * U[j] + U[j+1]
05D0
05D4
05D5 2001 DIVL2 -4(R7)[R5],R0 ; R0 <- R0 / V[1]
05DA 2002 MOVL R0,R3 ; Store quotient "digit" in R3
05DD 2003 BEQL 45$ ; Nothing to do if quotient is zero
05DF 2004 CMPL R3,#100 ; Is quotient LEQU 99?
05E1
05E6 2005 BLSSU 5$ ; Branch if quotient OK
05E8 2006 MOVL #99,R3 ; Otherwise start with 99
05EE
05EF 2007
05EF 2008 ; We will now multiply the divisor array by the quotient digit, storing the
05EF 2009 ; product in the scratch array.
05EF 2010
05EF 2011 5$: CLRL R2 ; Start out with a carry of zero
05F1 2012 MOVL R5,R4 ; R4 will be the loop counter
05F4 2013
05F4 2014 10$: EMUL R3,(R7)+,R2,R0 ; Multiply next divisor digit
05F8
05F9 2015 EDIV #100,R0,R2,(R8)+ ; Remainder to input array
05FF
0602 2016 ; Quotient becomes carry
0602 2017 SOBGTR R4,10$ ; More divisor digits?
0605 2018
0605 2019 MOVL R2,(R8)+ ; Store final carry
0608 2020
0608 2021 ; If the product array is larger than the dividend array, then the quotient is
0608 2022 ; too large. To avoid a second trip through the rather costly EMUL/EDIV loop,
0608 2023 ; and also to avoid array subtraction that produces a negative result, we will
0608 2024 ; first compare the product and dividend arrays. If the product is smaller, we
0608 2025 ; can safely subtract. If the product is larger, we decrease the quotient by
0608 2026 ; one and subtract the divisor array from the product array.
0608 2027
0608 2028 15$: MOVL R6,R0 ; Point R0 and R1 to high address ends
0608 2029 MOVL R8,R1 ; ... of dividend and scratch strings
060E 2030 MOVL R5,R4 ; Initialize the loop counter
0611 2031
0611 2032 ; The comparison is done from most to least significant digits
0611 2033
0611 2034 20$: CMPL -(R1),-(R0) ; Compare next pair of digits
0614 2035 BLSSU 30$ ; Leave loop if product is smaller
0616 2036 BGTRU 50$ ; Also leave if product is larger
0618 2037 SOBGEQ R4,20$ ; More to test?
061B 2038
061B 2039 ; If we drop through the loop, then the dividend and product are equal. We
061B 2040 ; simply store zeros in the dividend array (the equivalent of subtraction
061B 2041 ; of equal arrays) and return. Note that R0 is already pointing to the
061B 2042 ; least significant dividend array element.
061B 2043
061B 2044 MOVL R5,R4 ; Initialize still another loop counter
061E 2045

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```

      80      D4 061E 2046 25$: CLRL (R0)+ ; Store another zero
FB 54      F4 0620 2047      SOBGEQ R4,25$ ; Keep going?
          0623 2048
          05 0623 2049      RSB ; Return to caller
          0624 2050
          0624 2051 ; If we drop through the loop, then the quotient that is stored in R3 is good.
          0624 2052 ; We need to subtract the product array from the dividend array. Note that R0
          0624 2053 ; and R1 need to be adjusted to point to the least significant array elements
          0624 2054 ; before the subtraction can begin.
          0624 2055
          54 54 CE 0624 2056 30$ MNEGL R4,R4 ; We need a negative index
          50 6044 DE 0627 2057      MOVAL (R0)[R4],R0 ; Adjust dividend pointer
          51 6144 DE 0628 2058      MOVAL (R1)[R4],R1 ; ... and product pointer
          54 55 DO 062F 2059      MOVL R5,R4 ; R4 will count still another loop
          0632 2060
          80 81 C2 0632 2061 35$: SUBL2 (R1)+,(R0)+ ; Subtract next digits
          0A 18 0635 2062      BGEQ 40$ ; Skip to end of loop if no borrow
00000064 8F C0 0637 2063      ADDL2 #100,-4(R0) ; Add borrow back to this digit
          FC A0 063D 2064
          EE 54 D7 063F 2064      DECL (R0) ; ... and borrow from next highest digit
          F4 0641 2065 40$: SOBGEQ R4,35$ ; Keep going?
          0644 2066
          0644 2067 ; This is the exit path. R3 contains the quotient digit. The pointers to the
          0644 2068 ; various input and scratch arrays are in an indeterminate state.
          0644 2069
          05 0644 2070 45$: RSB ; Return to caller
          0645 2071
          0645 2072 ; The first guess at the quotient digit is too large. The brute force
          0645 2073 ; approach is to decrement the quotient by one and execute the EMUL/EDIV loop
          0645 2074 ; again. Note, however, that we can evaluate the modified product by
          0645 2075 ; subtracting the divisor from the initial product. Note also that, because
          0645 2076 ; the leading digit in the divisor is "large enough", we can only end up in
          0645 2077 ; this code path twice. (That is, the initial guess at the quotient will
          0645 2078 ; never be off by more than two.)
          0645 2079
          53 D7 0645 2080 50$: DECL R3 ; Try quotient smaller by one
          FB 13 0647 2081      BEQL 45$ ; All done if zero
          0649 2082
          0649 2083 ; Point R1 and R2 at the least significant digits of the scratch and product
          0649 2084 ; strings respectively.
          0649 2085
          50 55 CE 0649 2086      MNEGL R5,R0 ; Need a negative index
          51 FC A84C DE 064C 2087      MOVAL -4(R8)[R0],R1 ; Scratch array contains N+1 elements
          52 6740 DE 0651 2088      MOVAL (R7)[R0],R2 ; Product array contains N elements
          54 55 DO 0655 2089      MOVL R5,R4 ; R4 will count still another loop
          0658 2090
          81 82 C2 0658 2091 60$: SUBL2 (R2)+,(R1)+ ; Subtract next digits
          0A 18 0658 2092      BGEQ 70$ ; Skip to end of loop if no borrow
00000064 8F C0 065D 2093      ADDL2 #100,-4(R1) ; Add borrow back to this digit
          FC A1 0663 2094
          EE 54 D7 0665 2094      DECL (R1) ; ... and borrow from next highest digit
          F4 0667 2095 70$: SOBGEQ R4,60$ ; Keep going?
          066A 2096
          51 04 C0 066A 2097      ADDL2 #4,R1 ; Point R1 at most significant digit
          99 11 066D 2098      BRB 15$ ; Make another comparison
    
```

```

066F 2100 .SUBTITLE MULTIPLY_STRING - Multiply a String by a Number
066F 2101
066F 2102 : Functional Description:
066F 2103 :
066F 2104 : This routine multiplies an array of numbers (each array element LEQU
066F 2105 : 99) by a number (also LEQU 99). Each array element in the input array
066F 2106 : is replaced with the modified product, with the carry propogated to
066F 2107 : the next array element.
066F 2108
066F 2109 : Input Parameters:
066F 2110 :
066F 2111 : R3 - Multiplier
066F 2112 : R4 - Input array size
066F 2113 : R5 - Input array address
066F 2114
066F 2115 : Output Parameters:
066F 2116 :
066F 2117 : None
066F 2118
066F 2119 : Implicit Output:
066F 2120 :
066F 2121 : The input array elements are altered.
066F 2122
066F 2123 : Side Effects:
066F 2124 :
066F 2125 : R4 and R5 are modified by this routine.
066F 2126 :
066F 2127 : R3 is preserved.
066F 2128
066F 2129 : R0, R1, and R2 are used as scratch registers. R0 and R1 contain the
066F 2130 : quadword result of EMUL that is then passed into EDIV. R2 is the
066F 2131 : carry from one step to the next.
066F 2132
066F 2133 : Assumptions:
066F 2134 :
066F 2135 : This routine assumes that all array elements lie in the range from 0
066F 2136 : to 99 inclusive. (This is true if all input strings contain only legal
066F 2137 : decimal digits.) The arithmetic performed by this routine will
066F 2138 : maintain this assumption. The details of this argument can be found in
066F 2139 : the routine header for EXTENDED MULTIPLY_STRING. This routine performs
066F 2140 : less work so that those arguments also apply here.
066F 2141 :-
066F 2142

```

```

066F 2143 MULTIPLY_STRING:
52 D4 066F 2144 C'RL R2 ; Initial carry is zero
0671 2145
52 65 53 7A 0671 2146 10$: EMUL R3,(R5),R2,R0 ; Form modified product (R0 LEQU 9900)
0675 2147
00000064 8F 7B 0676 2147 EDIV #100,R0,R2,(R5)+ ; Remainder to input array
85 52 50 067C
067F 2148 ; Quotient becomes carry
EF 54 F5 067F 2149 SOBGTR R4,10$ ; Keep going?
0682 2150
65 52 D0 0682 2151 MOVL R2,(R5) ; Store final carry
05 0685 2152 RSB

```

E I M J E M M R C H - C O U S I E I C

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0686 2154 .SUBTITLE DECIMAL_ROPRAND
0686 2155 :-
0686 2156 : Functional Description:
0686 2157 :
0686 2158 : This routine receives control when a digit count larger than 31
0686 2159 : is detected. The exception is architecturally defined as an
0686 2160 : abort so there is no need to store intermediate state. All of the
0686 2161 : routines in this module save all registers R0 through R11 before
0686 2162 : performing the digit check. These registers must be restored
0686 2163 : before control is passed to VAX$ROPRAND.
0686 2164 :
0686 2165 : Input Parameters:
0686 2166 :
0686 2167 : 00(SP) - Saved R0
0686 2168 :
0686 2169 :
0686 2170 : 44(SP) - Saved R11
0686 2171 : 48(SP) - Return PC from VAX$xxxxxx routine
0686 2172 :
0686 2173 : Output Parameters:
0686 2174 :
0686 2175 : 00(SP) - Offset in packed register array to delta PC byte
0686 2176 : 04(SP) - Return PC from VAX$xxxxxx routine
0686 2177 :
0686 2178 : Implicit Output:
0686 2179 :
0626 2180 : This routine passes control to VAX$ROPRAND where further
0686 2181 : exception processing takes place.
0686 2182 :-
0686 2183 :
0686 2184 ASSUME ADDP6_B_DELTA_PC EQ ADDP4_B_DELTA_PC
0686 2185 ASSUME SUBP4_B_DELTA_PC EQ ADDP4_B_DELTA_PC
0686 2186 ASSUME SUBP6_B_DELTA_PC EQ ADDP4_B_DELTA_PC
0686 2187 ASSUME MULP_B_DELTA_PC EQ ADDP4_B_DELTA_PC
0686 2188 ASSUME DIVP_B_DELTA_PC EQ ADDP4_B_DELTA_PC
0686 2189 :
0686 2190 DECIMAL_ROPRAND:
OFFF 8F BA 0686 2191 POPR #M<R0,R1,R2,R3,R4,R5,R6,R7,R8,R9,R10,R11>
03 DD 068A 2192 PUSHL #ADDP4_B_DELTA_PC ; Store offset to delta PC byte
F971 31 068C 2193 BRW VAX$ROPRAND ; Pass control along

```

-S  
 ER  
 FC  
 LI  
 MT  
 CR  
 DE  
 LB

```

068F 2195      .SUBTITLE      ARITH_ACCVIO - Reflect an Access Violation
068F 2196      :+
068F 2197      : Functional Description:
068F 2198      :
068F 2199      : This routine receives control when an access violation occurs while
068F 2200      : executing within the emulator routines for ADDP4, ADDP6, SUBP4, SUBP6,
068F 2201      : MULP, or DIVP.
068F 2202      :
068F 2203      : The routine header for ASHP_ACCVIO in module VAX$ASHP contains a
068F 2204      : detailed description of access violation handling for the decimal
068F 2205      : string instructions.
068F 2206      :
068F 2207      : Input Parameters:
068F 2208      :
068F 2209      : See routine ASHP_ACCVIO in module VAX$ASHP
068F 2210      :
068F 2211      : Output Parameters:
068F 2212      :
068F 2213      : See routine ASHP_ACCVIO in module VAX$ASHP
068F 2214      :-
068F 2215
068F 2216      ARITH_ACCVIO:
068F 2217      CLRL      R2          ; Initialize the counter
068F 2218      PUSHAB  MODULE_BASE ; Store base address of this module
068F 2219      PUSHAB  MODULE_END   ; Store module end address
068F 2220      BSBW    DECIMAL$BOUNDS_CHECK ; Check if PC is inside the module
068F 2221      ADDL   #4,SP         ; Discard end address
068F 2222      SUBL2  (SP)+,R1      ; Get PC relative to this base
068F 2223
068F 2224      10$:  CMPW    R1,PC_TABLE_BASE[R2] ; Is this the right PC?
068F 2225      BEQL   30$         ; Exit loop if true
068F 2226      AOBLS  #TABLE_SIZE,R2,10$ ; Do the entire table
068F 2227
068F 2228      ; If we drop through the dispatching based on PC, then the exception is not
068F 2229      ; one that we want to back up. We simply reflect the exception to the user.
068F 2230
068F 2231      20$:  POPR    #*M<R0,R1,R2,R3> ; Restore saved registers
068F 2232      RSB     ; Return to exception dispatcher
068F 2233
068F 2234      ; The exception PC matched one of the entries in our PC table. R2 contains
068F 2235      ; the index into both the PC table and the handler table. R1 has served
068F 2236      ; its purpose and can be used as a scratch register.
068F 2237
068F 2238      30$:  MOVZWL  HANDLER_TABLE_BASE[R2],R1 ; Get the offset to the handler
068F 2239      JMP     MODULE_BASE[R1] ; Pass control to the handler
068F 2240
068F 2241      ; In all of the instruction-specific routines, the state of the stack
068F 2242      ; will be shown as it was when the exception occurred. All offsets will
068F 2243      ; be pictured relative to R0.

```

```

          52      D4
F96B CF      9F
06FC'CF      9F
          F964'  30
          5E      04      C0
          51      8E      C2
0000'CF42    51      B1
          07      13
          F4 52  2A      F2
          OF      BA
          05
51 0000'CF42  3C
          F944 CF41  17

```

```

06BC 2245      .SUBTITLE      Access Violation Handling for ADDPx and SUBPx
06BC 2246      :+
06BC 2247      : Functional Description.
06BC 2248      :
06BC 2249      : The only difference among the various entry points is the number of
06BC 2250      : longwords on the stack. R0 is advanced beyond these longwords to point
06BC 2251      : to the list of saved registers. These registers are then restored,
06BC 2252      : effectively backing the routine up to its initial state.
06BC 2253      :
06BC 2254      : Input Parameters:
06BC 2255      :
06BC 2256      : R0 - Address of top of stack when access violation occurred
06BC 2257      :
06BC 2258      : See specific entry points for details
06BC 2259      :
06BC 2260      : Output Parameters:
06BC 2261      :
06BC 2262      : See input parameter list for VAX$DECIMAL_ACCVIO in module VAX$ASHP
06BC 2263      :-
06BC 2264      :
06BC 2265      :+
06BC 2266      : ADD_SUB_BSBW_24
06BC 2267      :
06BC 2268      : An access violation occurred in one of the subroutines ADD_PACKED_BYTE,
06BC 2269      : SUB_PACKED_BYTE, or STORE_RESULT. In addition to the six longwords of work
06BC 2270      : space, this routine has an additional longword, the return PC, on the
06BC 2271      : stack.
06BC 2272      :
06BC 2273      : 00(R0) - Return PC in mainline VAX$xxxxxx routine
06BC 2274      : 04(R0) - Address of sign byte of destination string
06BC 2275      : 08(R0) - First longword of scratch space
06BC 2276      : etc.
06BC 2277      :-
06BC 2278      :
50   04   C0 06BC 2279 ADD_SUB_BSBW_24:
06BC 2280      ADDL    #4,R0           ; Skip over return PC and drop into ...
06BF 2281      :
06BF 2282      :+
06BF 2283      : ADD_SUB_24
06BF 2284      :
06BF 2285      : There are five longwords of workspace and a saved string address on the stack
06BF 2286      : for this entry point.
06BF 2287      :
06BF 2288      : 00(R0) - Address of sign byte of destination string
06BF 2289      : 04(R0) - First longword of scratch space
06BF 2290      :
06BF 2291      :
06BF 2292      : 20(R0) - Fifth longword of scratch space
06BF 2293      : 24(SP) - Saved R0
06BF 2294      : 28(SP) - Saved R1
06BF 2295      : etc.
06BF 2296      :-
06BF 2297      :
50   18   C0 06BF 2298 ADD_SUB_24:
   F93B'  31 06BF 2299      ADDL    #24,R0          ; Discard scratch space on stack
06C2 2300      BRW     VAX$DECIMAL_ACCVIO  ; Join common code to restore registers
06C5 2301
  
```

Sy  
 --  
 \$U  
 BE  
 BI  
 BR  
 BR  
 BU  
 BU  
 BU  
 BU  
 CL  
 CL  
 CL  
 CL  
 CL  
 CL  
 CL  
 CL  
 DC  
 DE  
 DE  
 DE  
 DE  
 DE  
 DE  
 DE  
 DE  
 DI  
 DI  
 DI  
 DI  
 DU  
 ER  
 ER  
 ER  
 ER













VAX\$DECIMAL\_ARITHMETIC  
Symbol table

```

...PC... = 00000553
...ROPRAND = 000004BF R 02
ADDP4_B_DELTA_PC = 00000003
ADDP6_B_DELTA_PC = 00000003
ADD_PACKED = 0C0000E5 R 02
ADD_PACKED_BYTE_R6_R7 = 00000170 R 02
ADD_PACKED_BYTE_STRING = 0000016A R 02
ADD_SUBTRACT_EXIT = 0000013D R 02
ADD_SUB_24 = 000006BF R 02
ADD_SUB_BSBW_0 = 000006C8 R 02
ADD_SUB_BSBW_24 = 000006BC R 02
ADD_SUB_BSBW_4 = 000006C5 R 02
ADD_SUB_V_ZERO_R4 = = 0000001F
ARITH_ACCVIO = 0000068F R 02
CHECK_WRITE_ACCESS = 00000292 R 02
DECIMAL$BINARY_TO_PACKED_TABLE ***** X 00
DECIMAL$BOUNDS_CHECK ***** X 00
DECIMAL$PACKED_TO_BINARY_TABLE ***** X 00
DECIMAL$STRIP_ZEROS_R0_R1 ***** X 00
DECIMAL$STRIP_ZEROS_R2_R3 ***** X 00
DECIMAL_ROPRAND = 00000686 R 02
DIVIDE_BY_ZERO = 0000049E R 02
DIVP_0 = 000006F1 R 02
DIVP_BSBW_0 = 000006EE R 02
DIVP_B_DELTA_PC = = 00000003
DIVP_R6_R7 = 000006F4 R 02
EXTEND_STRING_MULTIPLY = 00000470 R 02
HANDLER_TABLE_BASE = 00000000 R 04
MODULE_BASE = 00000000 R 02
MODULE_END = = 000006FC R 02
MULP_AT_SP = 000006D4 R 02
MULP_BSBW_0 = 000006EE R 02
MULP_B_DELTA_PC = = 00000003
MULP_DIVP_0 = 000006F1 R 02
MULP_DIVP_8 = 000006E8 R 02
MULP_DIVP_R9 = 000006E0 R 02
MULP_R8 = 000006CE R 02
MULTIPLY_DIVIDE_EXIT = 0000037D R 02
MULTIPLY_STRING = 0000066F R 02
PC_TABLE_BASE = 00000000 R 03
PSL$M_N = = 00000008
PSL$M_V = = 00000002
PSL$M_Z = = 00000004
PSL$V_CM = = 0000001F
PSL$V_V = = 00000001
PSL$V_Z = = 0C000002
QUOTIENT_DIGIT = 000005CA R 02
SRMSK_FLY_DIV_T = = 00000004
STORE_RESULT = 00000254 R 02
SUBP4_B_DELTA_PC = = 00000003
SUBP6_B_DELTA_PC = = 00000003
SUBTRACT_PACKED = 00000198 R 02
SUB_PACKED_BYTE_R6_R7 = 0000022E R 02
SUB_PACKED_BYTE_STRING = 00000228 R 02
TABLE_SIZE = 0000002A
VAX$ADDP4 = 00000030 RG 02
VAX$ADDP6 = 00000009 RG 02

```

```

VAX$ADD_PACKED_BYTE_R6_R7 00000170 RG 02
VAX$DECIMAL_ACCVIO ***** X 00
VAX$DECIMAL_EXIT ***** X 00
VAX$DIVP 000004B1 RG 02
VAX$MULP 000002AD RG 02
VAX$REFLECT_TRAP ***** X 00
VAX$ROPRAND ***** X 00
VAX$SUBP4 00000022 RG 02
VAX$SUBP6 00000000 RG 02

```

-----  
! Psect synopsis !  
-----

PSECT name	Allocation	PSECT No.	Attributes
ABS	00000000 ( 0.)	00 ( 0.)	NOPIC USR CON ABS LCL NOSHR NOEXE NORD NOWRT NOVEC BYTE
\$ABSS	00000000 ( 0.)	01 ( 1.)	NOPIC USR CON ABS LCL NOSHR EXE RD WRT NOVEC BYTE
VAX\$CODE	000006FC ( 1788.)	02 ( 2.)	PIC USR CON REL LCL SHR EXE RD NOWRT NOVEC LONG
PC TABLE	00000054 ( 84.)	03 ( 3.)	PIC USR CON REL LCL SHR NOEXE RD NOWRT NOVEC BYTE
HANDLER_TABLE	00000054 ( 84.)	04 ( 4.)	PIC USR CON REL LCL SHR NOEXE RD NOWRT NOVEC BYTE

-----  
! Performance indicators !  
-----

Phase	Page faults	CPU Time	Elapsed Time
Initialization	75	00:00:00.20	00:00:02.24
Command processing	79	00:00:00.46	00:00:03.73
Pass 1	241	00:00:09.83	00:00:40.27
Symbol table sort	0	00:00:00.42	00:00:02.26
Pass 2	393	00:00:06.36	00:00:32.16
Symbol table output	0	00:00:00.08	00:00:00.54
Psect synopsis output	0	00:00:00.04	00:00:00.36
Cross-reference output	0	00:00:00.00	00:00:00.00
Assembler run totals	788	00:00:17.39	00:01:21.56

The working set limit was 1800 pages.  
52274 bytes (103 pages) of virtual memory were used to buffer the intermediate code.  
There were 20 pages of symbol table space allocated to hold 184 non-local and 116 local symbols.  
2646 source lines were read in Pass 1, producing 25 object records in Pass 2.  
23 pages of virtual memory were used to define 21 macros.

-----  
! Macro library statistics !  
-----

Macro library name	Macros defined
-\$255\$DUA18:[EMULAT.OBJ]VAXMACROS.MLB;1	12
-\$255\$DUA18:[SYSLIB]STARLET.MLB;3	6
TOTALS (all libraries)	18

318 GETS were required to define 18 macros.

There were no errors, warnings or information messages.

MACRO/LIS=LIS\$:VAXARITH/OBJ=OBJ\$:VAXARITH MSRC\$:VAXARITH/UPDATE=(BUG\$:VAXARITH)+LIB\$:VAXMACROS/LIB

0441 AH-EF71A-SE  
VAX/VMS V4.1 SRC LST MCRF UPD

A dense grid of source code listings for various VAX/VMS components. The grid is organized into several major sections, each containing multiple columns of code. The sections are labeled as follows:

- VAXARITH LIS**: Located in the upper-middle section, containing arithmetic-related source code.
- EMULAT**: Located in the middle section, containing emulation-related source code.
- VAXEMUL MAP**: Located in the middle section, containing mapping-related source code.
- ERFPROC1 MAP**: Located in the middle-right section, containing process-related source code.
- ERF**: Located in the lower-middle section, containing error-related source code.
- ERF MAP**: Located in the lower-middle section, containing error-related source code.
- RESELECT LIS**: Located in the lower-right section, containing selection-related source code.

The code listings are presented in a standard VAX/VMS source listing format, with line numbers and column markers visible on the left side of each listing block.